

Advanced Network Programming

Software Defined Networking

Lin Wang
Fall 2020, Period 1



Part of the content is adapted from Scott Shenker

Part 2: network infrastructure

Lecture 7: Network forwarding and routing

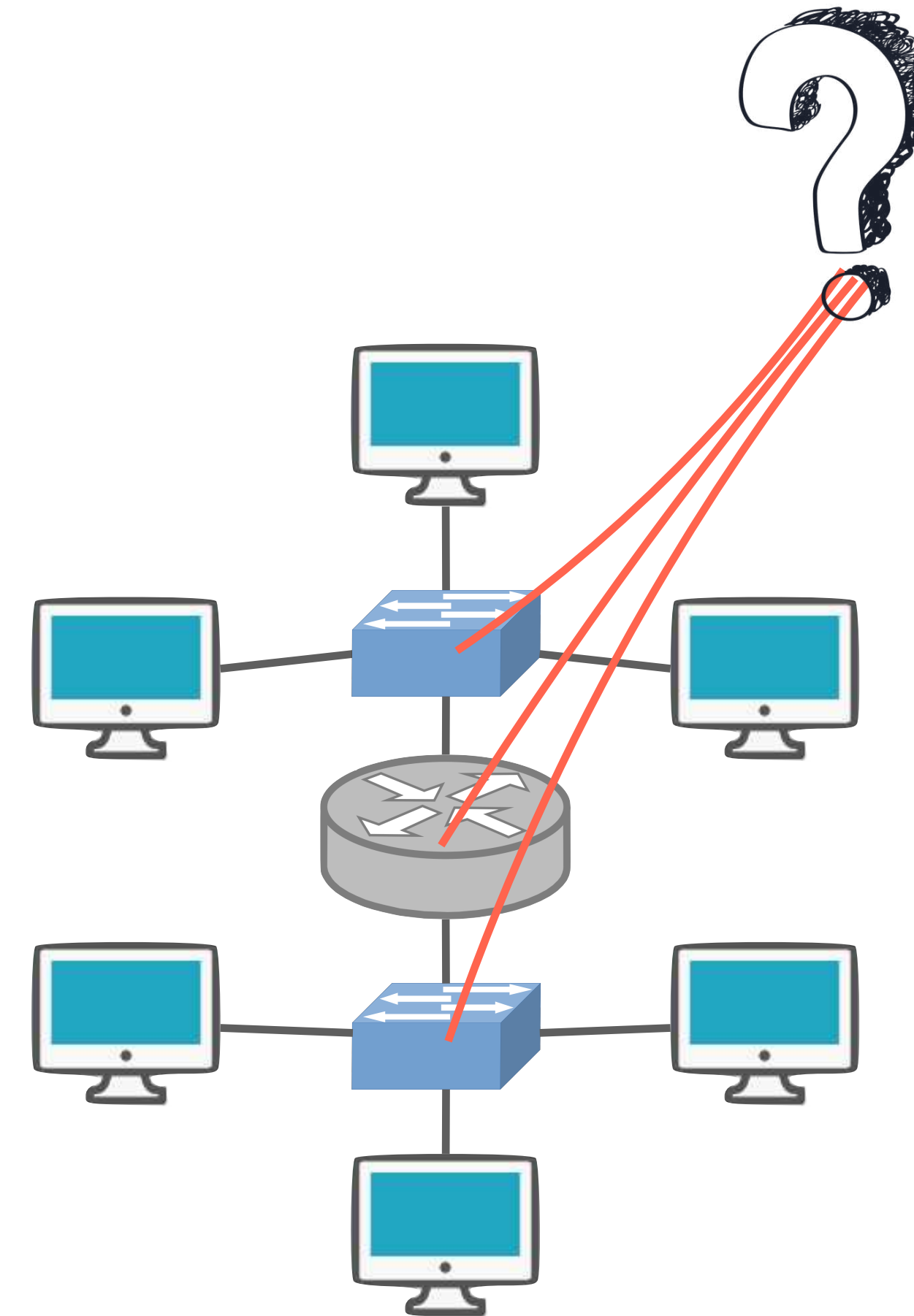
Lecture 8: Software defined networking

- Complexity in network control and management
- Centralized control, OpenFlow
- Network virtualization, FlowVisor

Lecture 9: Programmable data plane

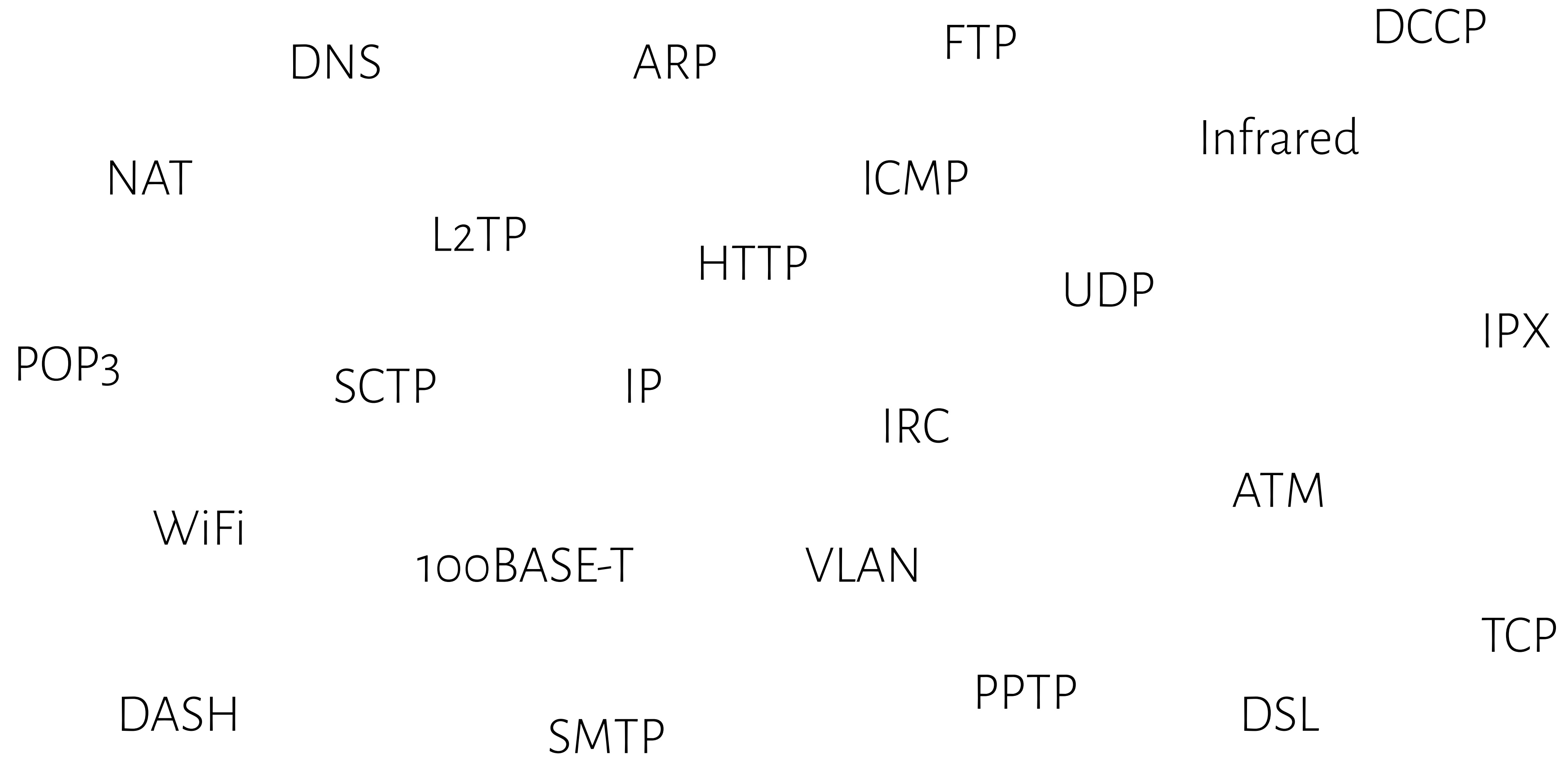
Lecture 10: Cloud networking

Lecture 11: Beyond networking



**What is your impression about computer
networking so far?**

A plethora of protocol acronyms



A word cloud of network protocol acronyms. The acronyms are arranged in a roughly circular pattern, with some appearing more frequently or in larger fonts than others. The acronyms include: DNS, ARP, FTP, DCCP, NAT, ICMP, Infrared, L2TP, HTTP, UDP, POP₃, SCTP, IP, IRC, ATM, WiFi, 100BASE-T, VLAN, TCP, DASH, SMTP, PPTP, and DSL.

DNS

ARP

FTP

DCCP

NAT

ICMP

Infrared

L2TP

HTTP

UDP

POP₃

SCTP

IP

IRC

ATM

WiFi

100BASE-T

VLAN

TCP

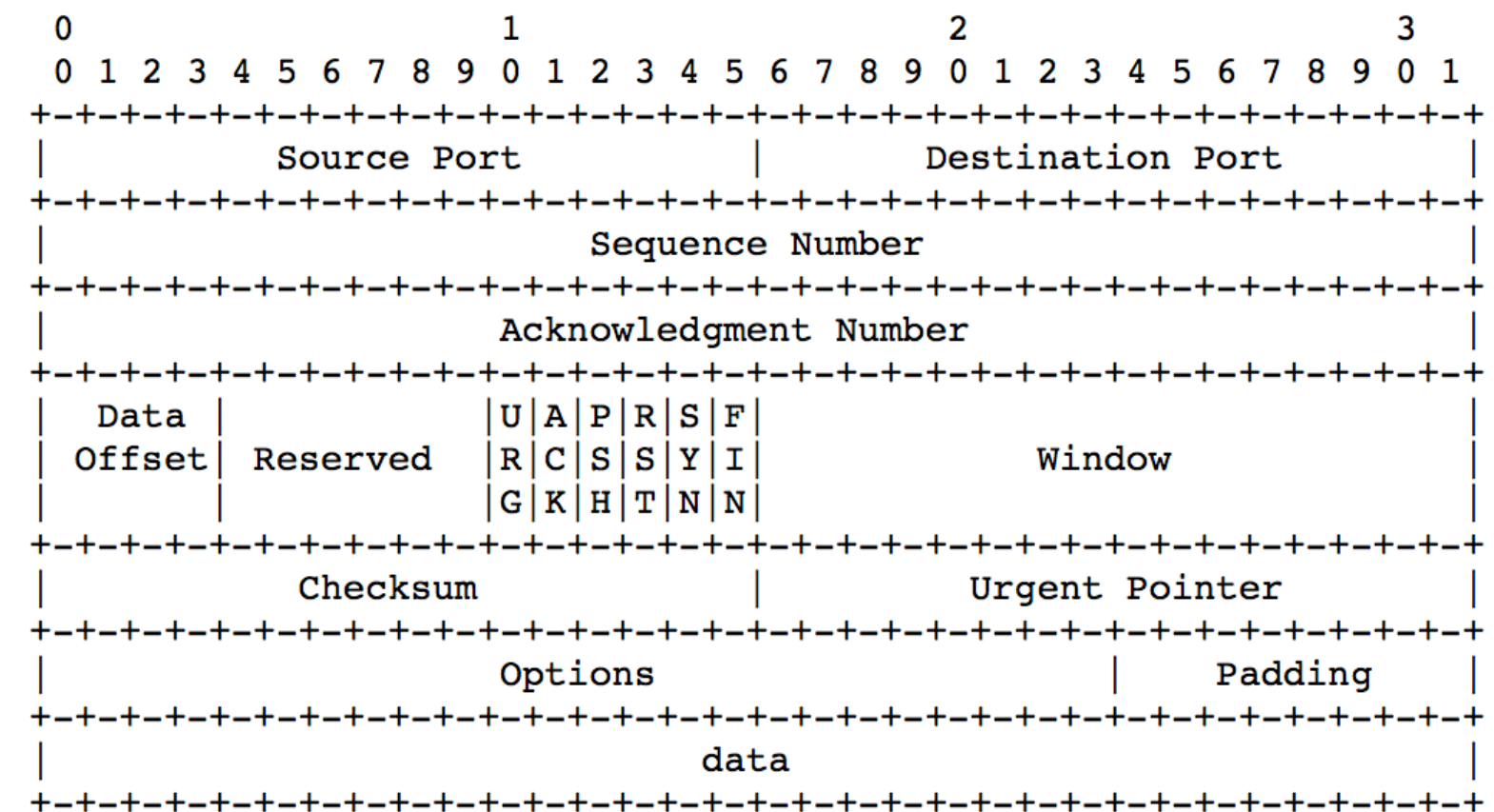
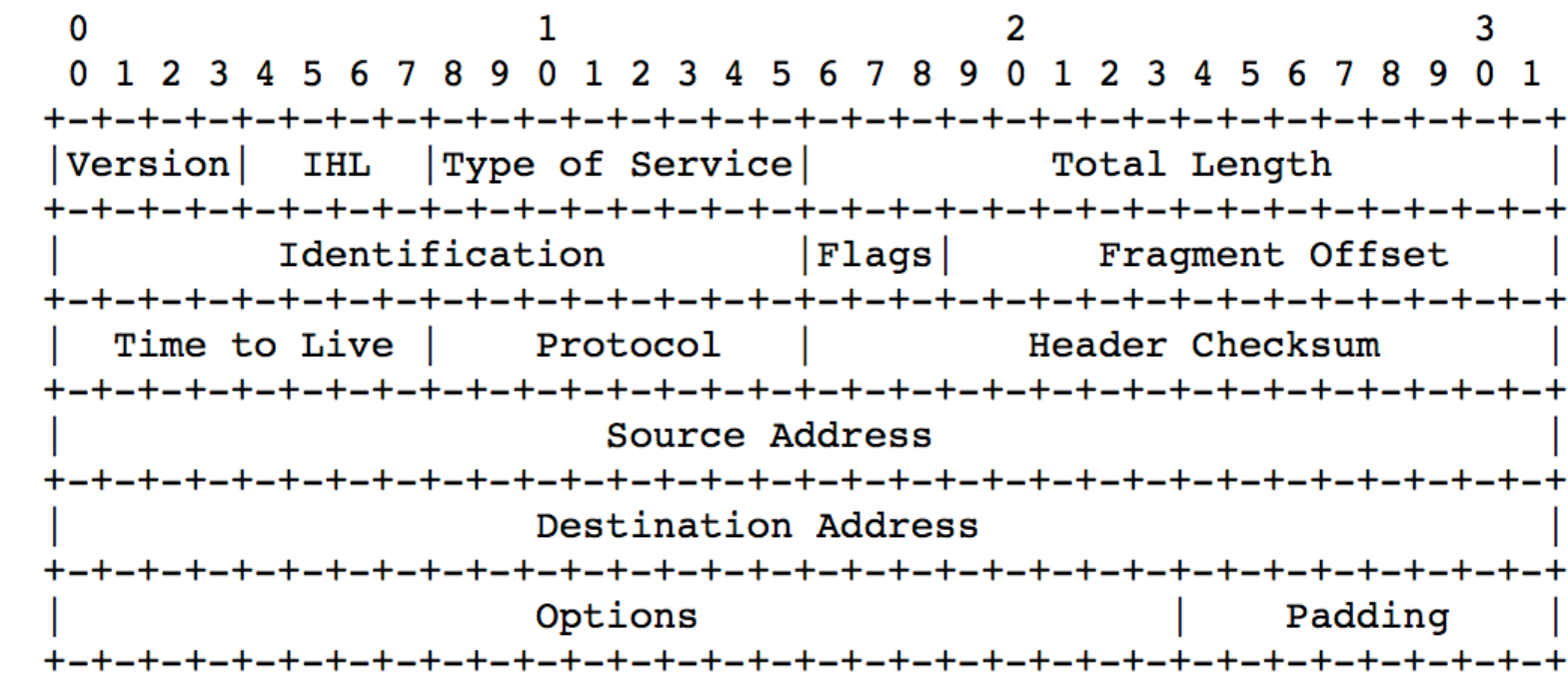
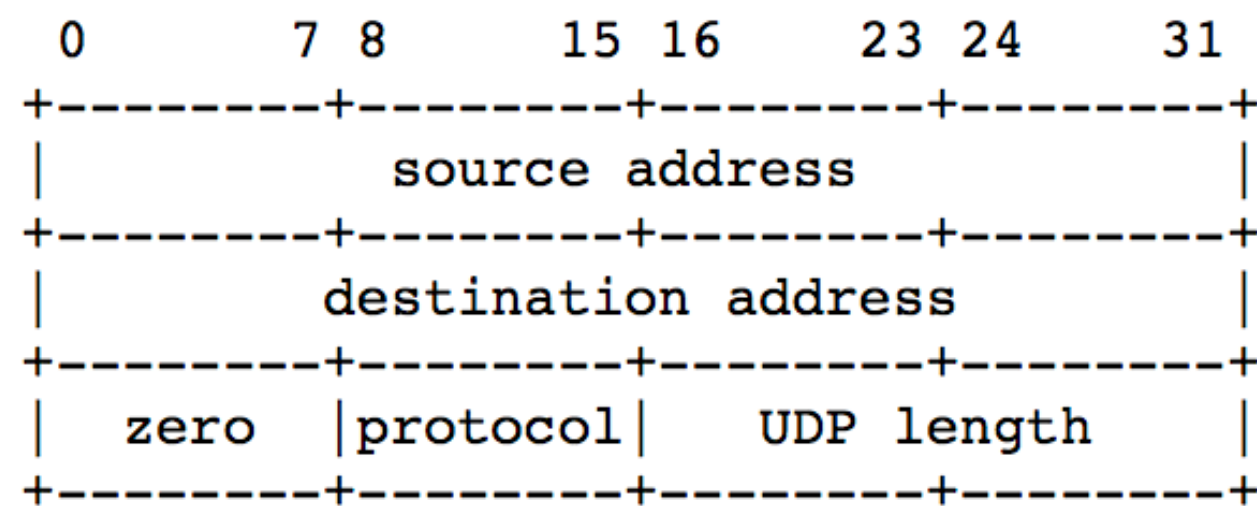
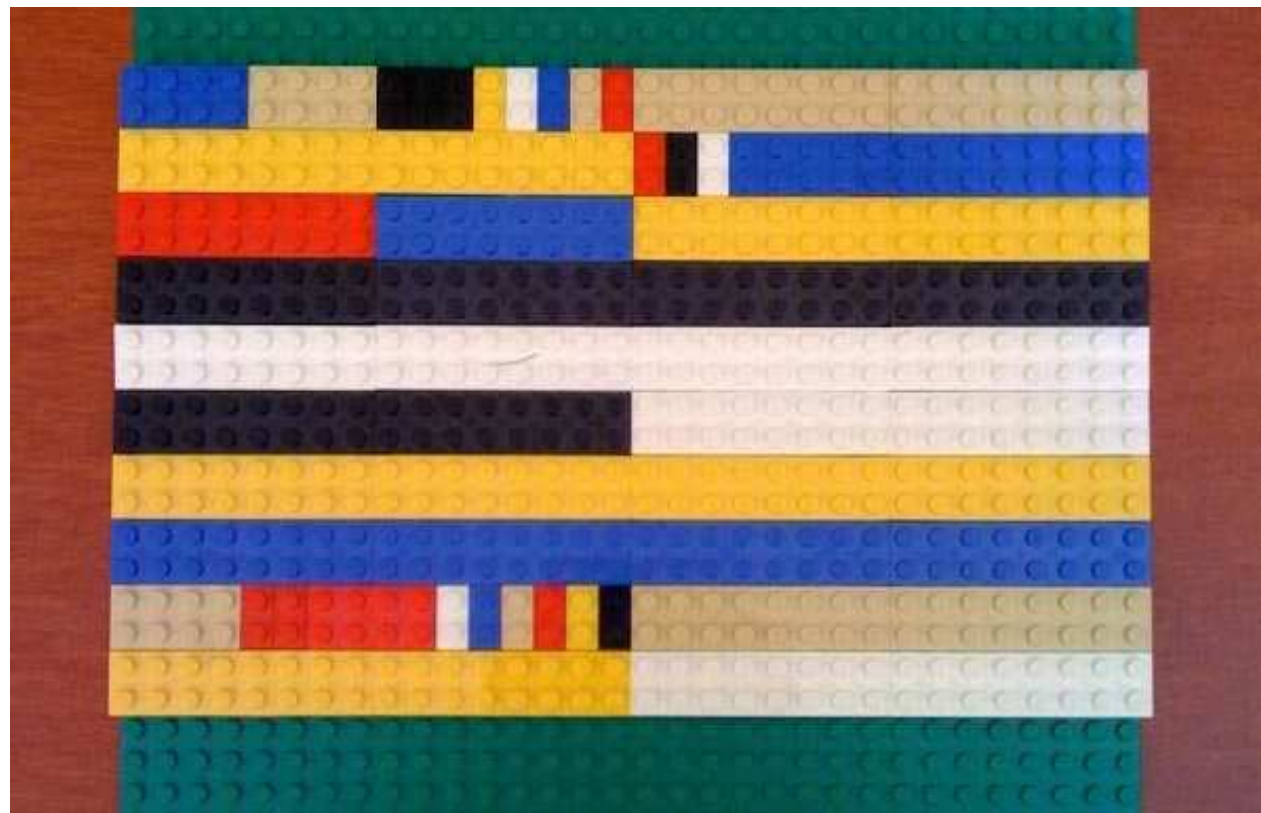
DASH

SMTP

PPTP

DSL

A heap of header formats



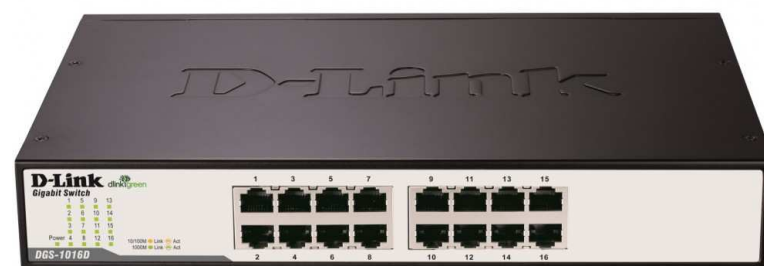
A bunch of boxes



Access Point



Router



Switch



Middlebox (Firewall, Load
Balancer, IDS)

A ton of tools

A collection of network tools arranged in a circular pattern. The tools are: ping, tcpdump, wget, nmap, syslog, ntop, whois, dig, nslookup, rancid, iperf, traceroute, and wireshark. The tools are arranged in a circular pattern, with ping at the top and whois at the bottom.

ping

tcpdump

wget

nmap

syslog

ntop

whois

dig

nslookup

rancid

iperf

traceroute

wireshark

Why are there so many artifacts?

Complexity in networking

We need **different functionalities**, also new ones

- Different physical layers and applications, traffic engineering, congestion control, security

Networks run in a **distributed, autonomous** way

- Scalability is important

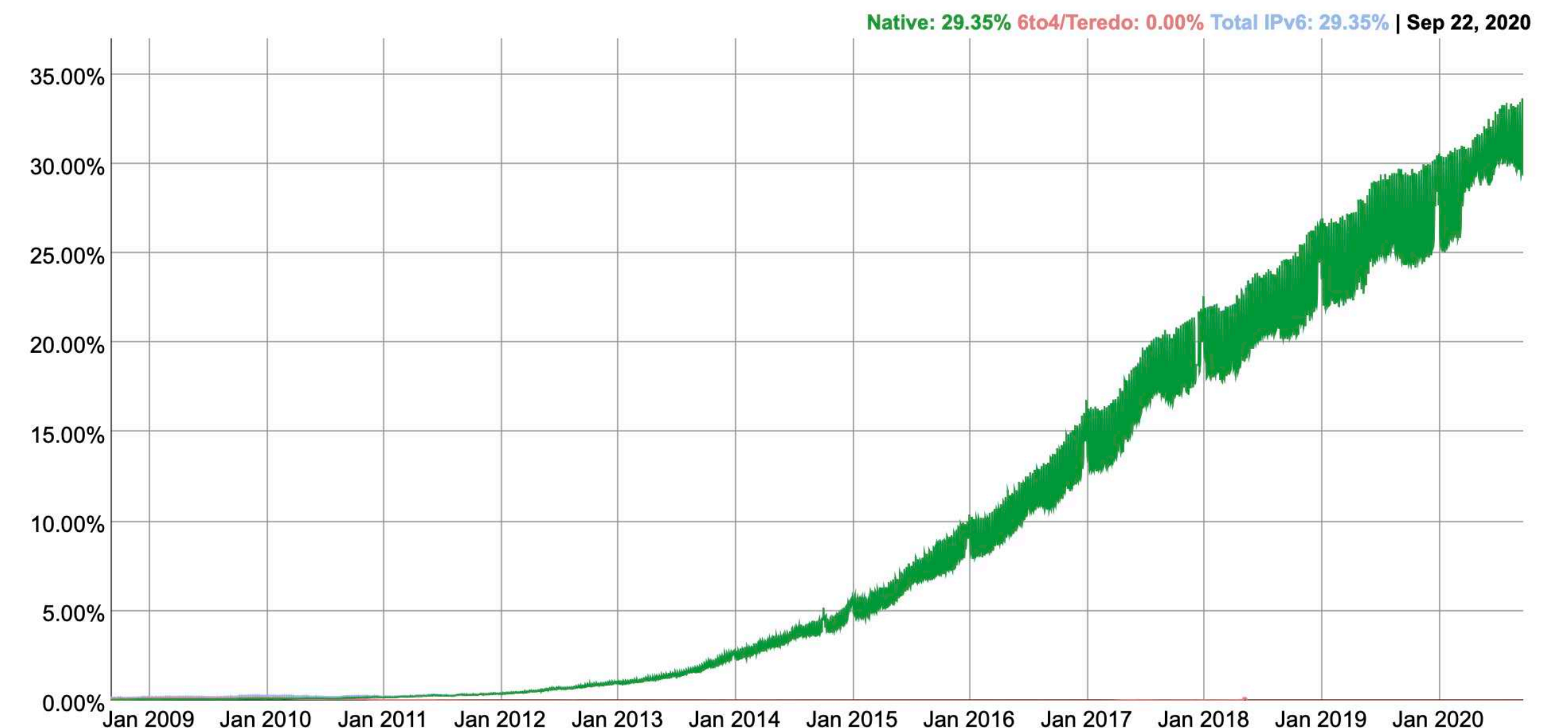
All these add to **complexity**, innovations are active in academia, but suffer from poor adoption of deployment

- Example: IPv6
- Deadlock between innovation and adoption

RFC 2460 (1998)

IPv6 Adoption

We are continuously measuring the availability of IPv6 connectivity among Google users. The graph shows the percentage of users that access Google over IPv6.



<https://www.google.com/intl/en/ipv6/statistics.html#tab=ipv6-adoption>

Once upon a time

AT&T updates their internal network infrastructure (routers and switches) **every 18 months** to keep up with the current demands for network



One upgrade requires **billion of USD**

- A Cisco top of line switch costs \$27K USD
- Significant **manpower** is needed to upgrade the network

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AT&T Targets Flexibility, Cost Savings With New Network Design

Move Could Cut the Company's Capital Costs by Billions of Dollars

By [Thomas Gryta](#)
Updated Feb. 24, 2014 12:25 pm ET

PRINT TEXT

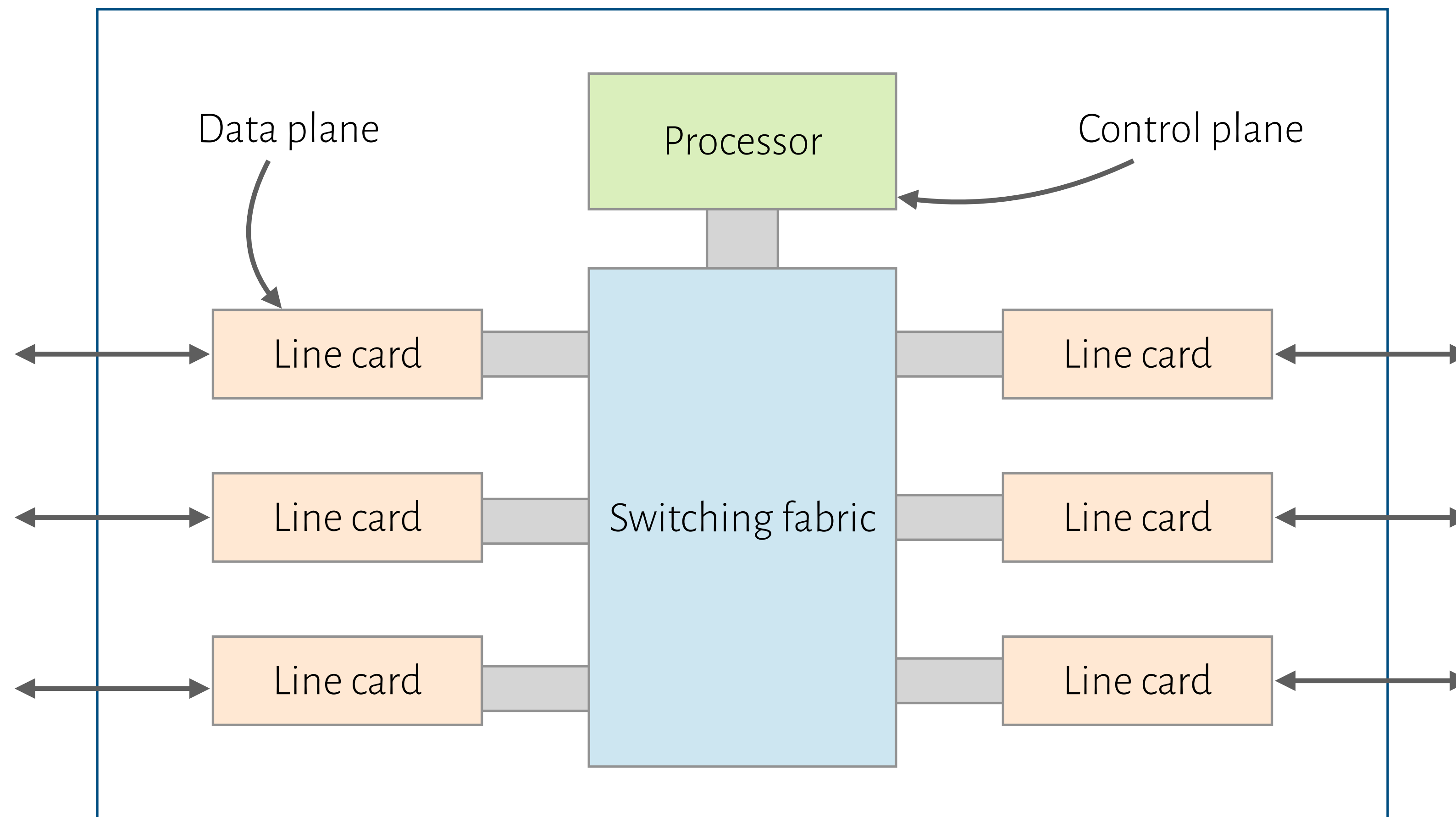
AT&T Inc. is planning to rebuild its sprawling network with less expensive, off-the-shelf equipment controlled by software, a move that could cut its capital costs by billions of dollars and put further pressure on telecom gear makers.

The shift will mean the second-largest U.S. carrier will buy less specialized equipment from vendors such as Ericsson, Alcatel-Lucent SA and Cisco Systems Inc., and instead purchase more generic hardware from a wider variety of producers. That equipment will be tied together with software, making it easier and cheaper to upgrade to new technologies, roll out new services or respond to changes in demand for connectivity.

RECOMMENDED VIDEOS

1. Biden Takes Aim at Trump on Social Security
2. Gyms Brace for Fewer Members as At-Home Fitness Rises During Pandemic
3. Some Covid-19 Survivors Grapple With Large Medical Bills
4. The Rise of TikTok: From China to America

Complexity in network planes



Complexity in the control plane

Control plane needs to achieve goals such as connectivity, inter-domain policy, isolation, access control...

Currently, these goals are achieved by many mechanisms/protocols:

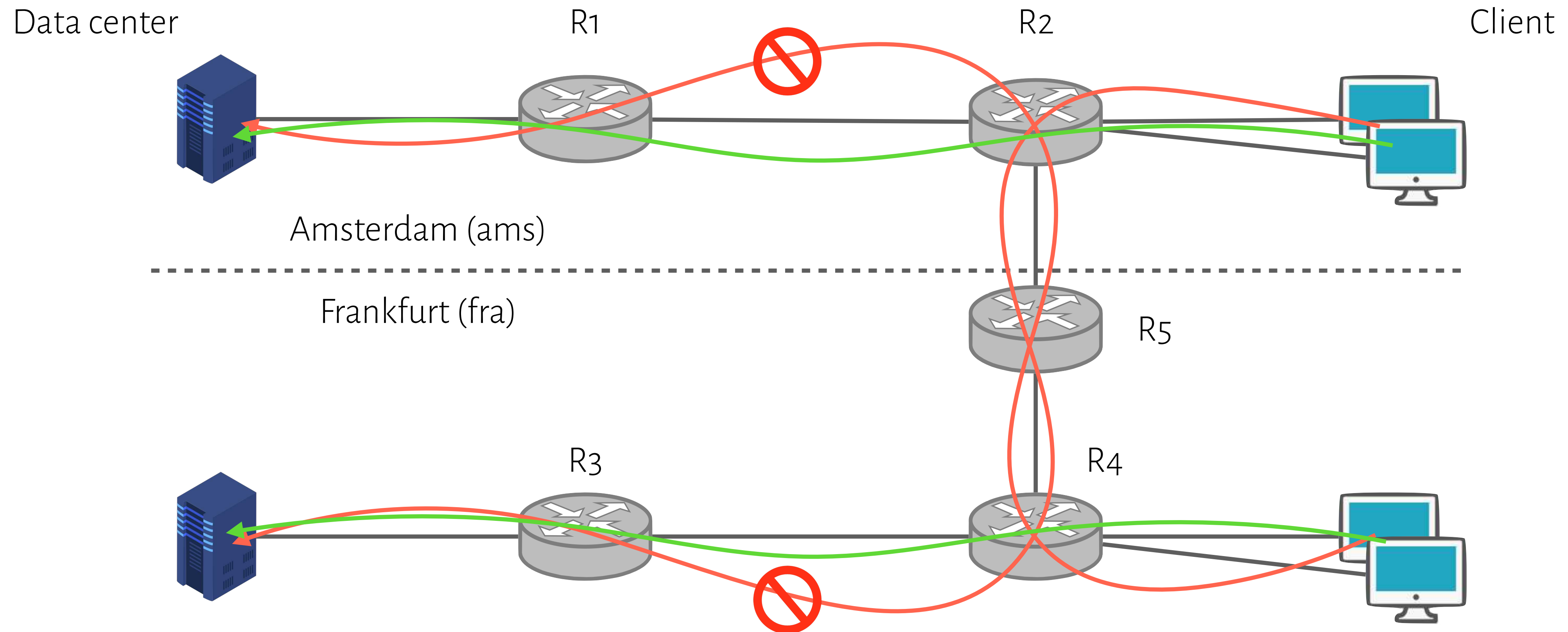
- Globally **distributed**: routing algorithms
- **Manual**/scripted configuration: Access Control Lists, VLANs
- **Centralized** computation: traffic engineering (indirect control)

Even worse, these mechanisms/protocols **interact with each other**

- Routing, addressing, access control, QoS

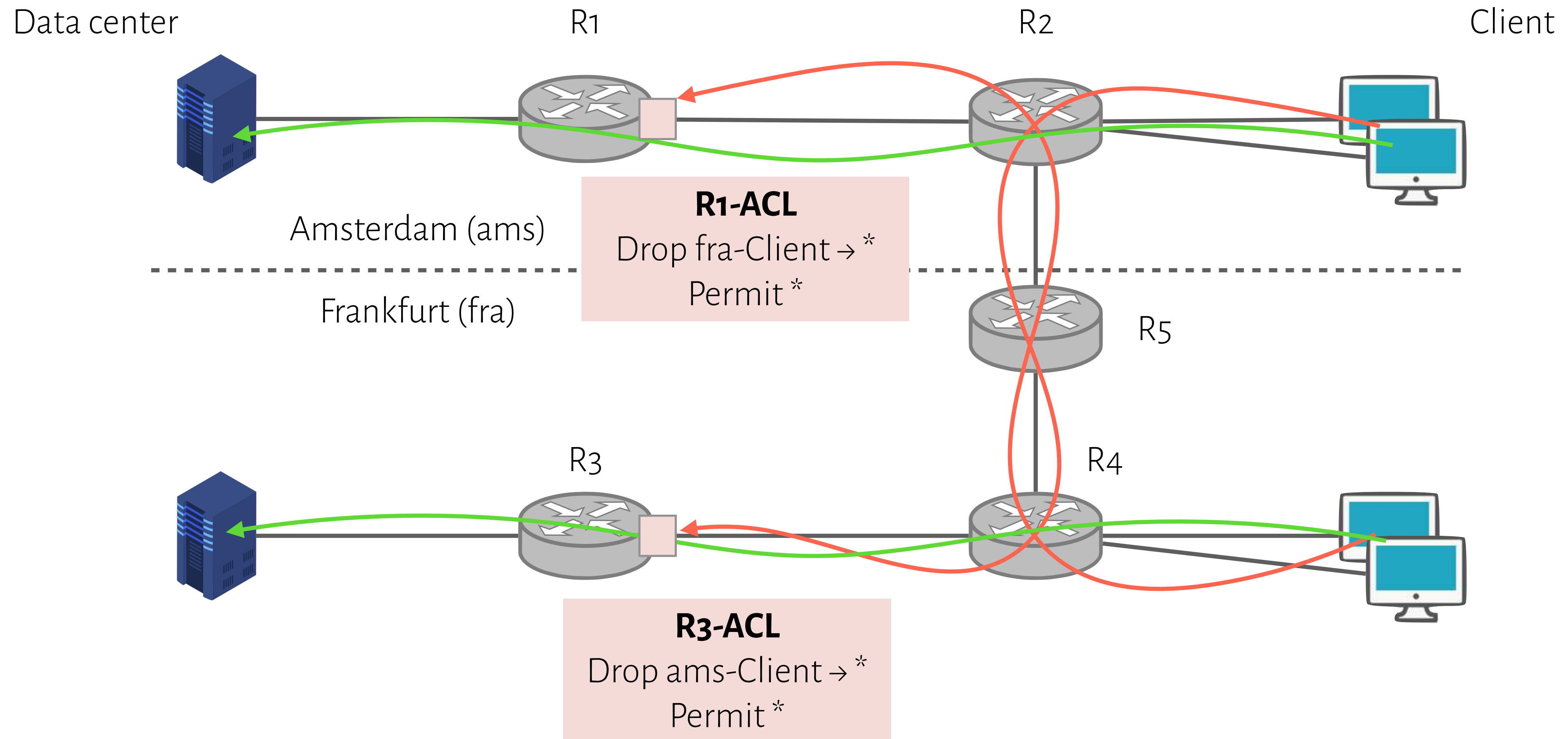
Network control plane is a complicated mess!

Example: access control configuration

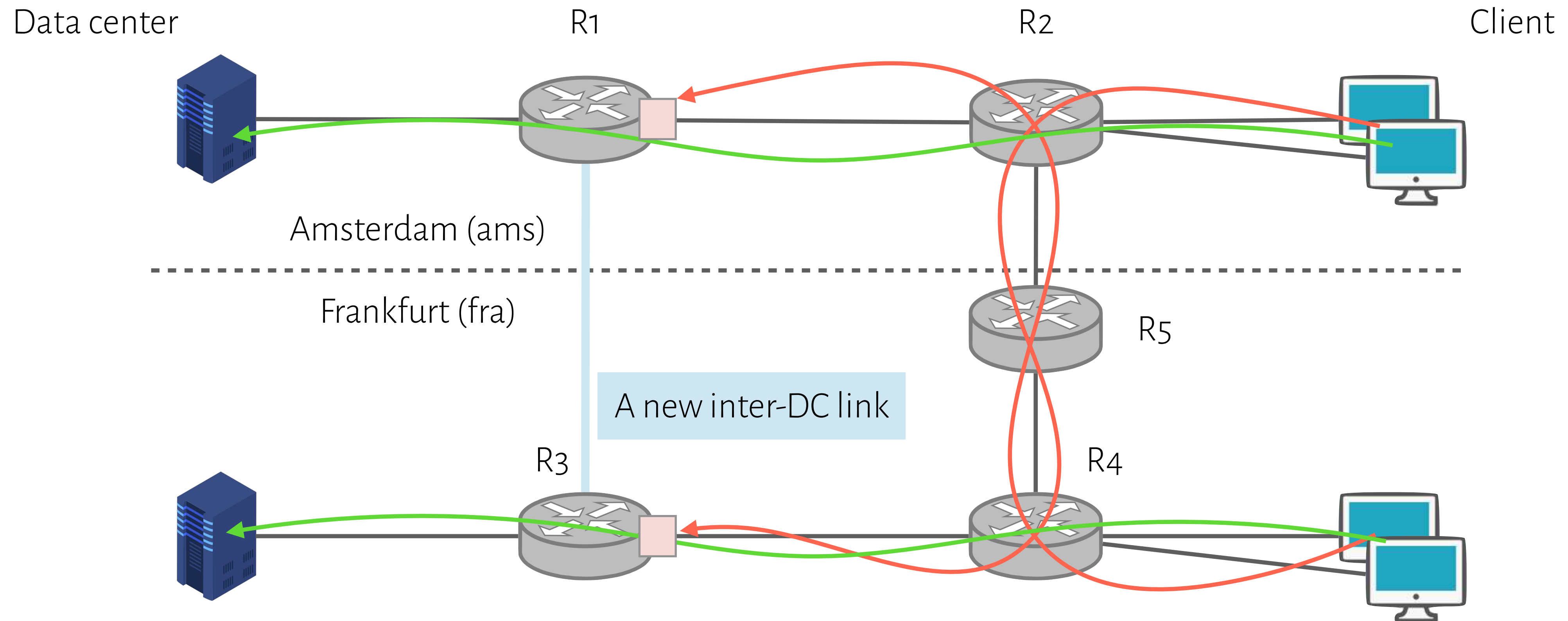


How can we block cross-region traffic?

Example: access control configuration

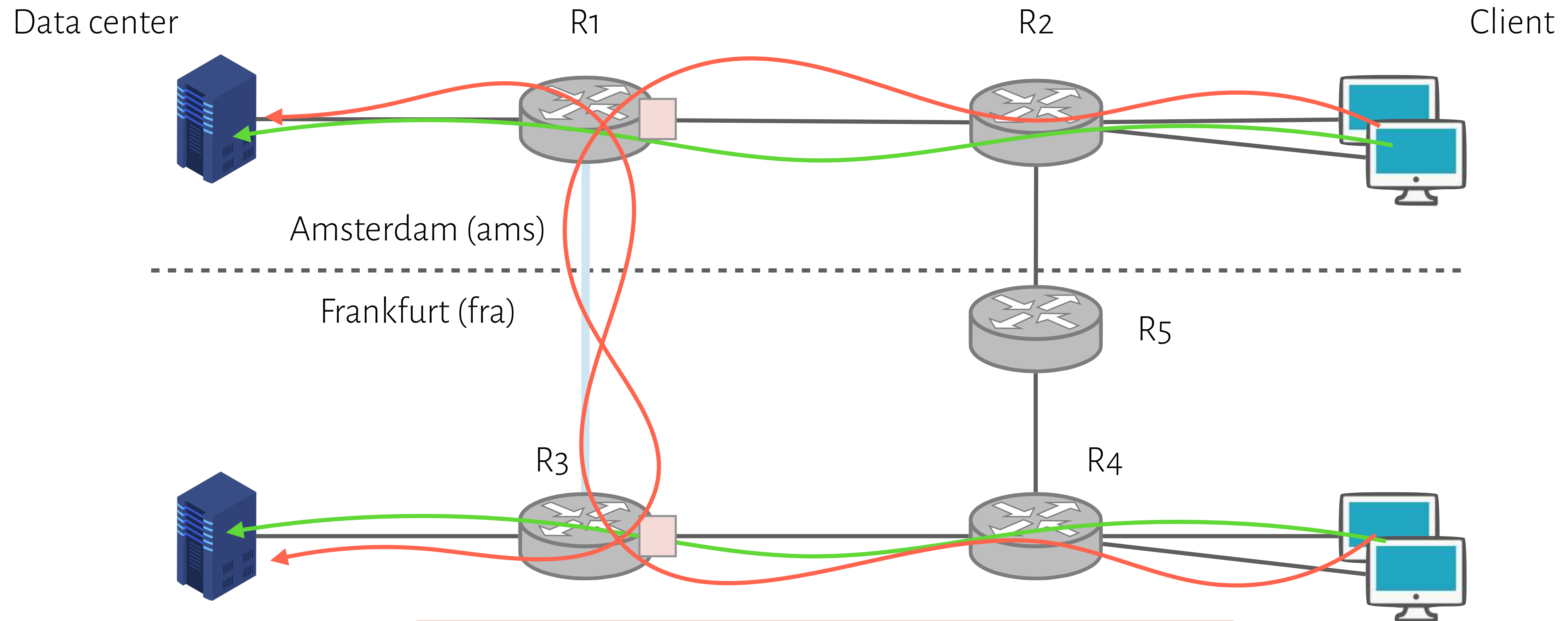


Example: access control configuration



What would be the problem?

Example: access control configuration



Network traffic can bypass the ACLs now.
→ Routing and access control are dependent!

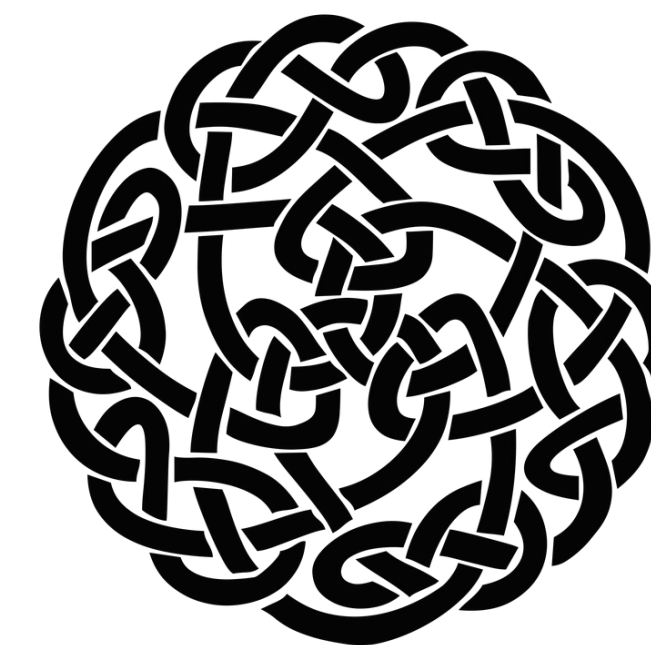
How have we managed to survive?

Network administrators miraculously master this complexity

- Understand all aspects of networks
- Must keep myriad details in mind

The ability to master complexity is both a blessing and a curse!

The **ability to master complexity** is valuable but not the same as **the ability to extract simplicity**



The Gordian Knot

Actually, networks work better during the weekends

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Fewer heart attack patients die when top cardiologists are away at conferences, study finds

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Specialists may do more harm than good, a study suggests CREDIT: LYNNE CAMERON PA

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By **Sarah Knapton**, SCIENCE EDITOR
9 MARCH 2018 • 2:56PM

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Monday

Tuesday

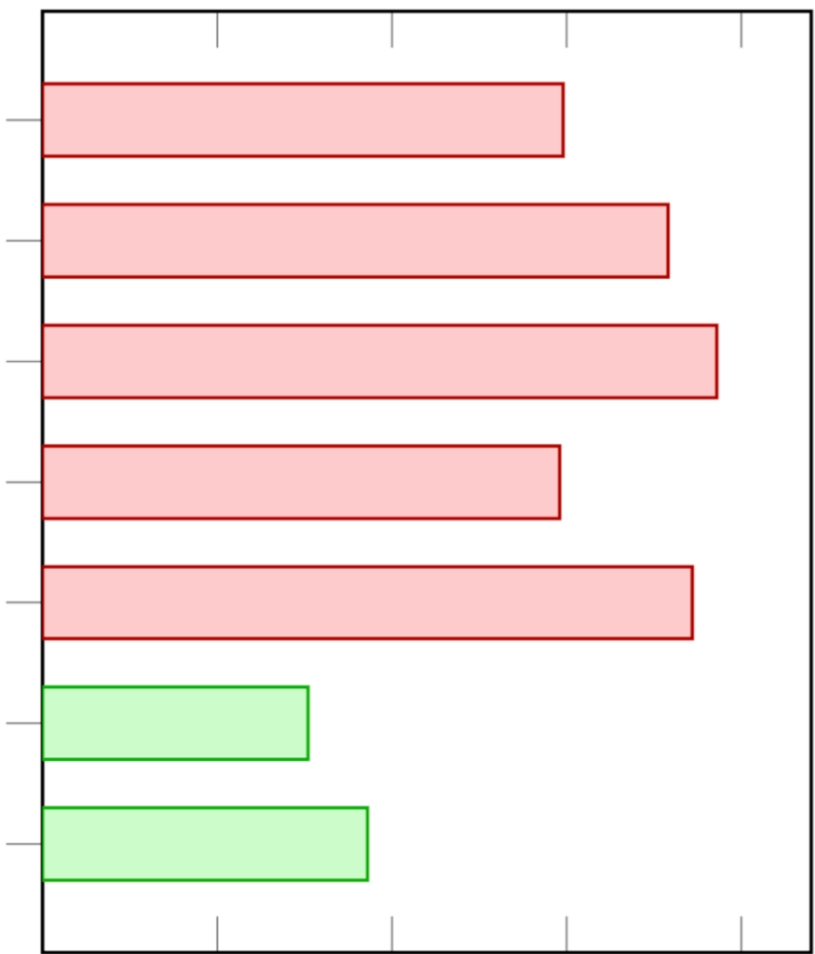
Wednesday

Thursday

Friday

Saturday

Sunday



0 5 10 15 20

% of route leaks

source: Job Snijders (NTT)

How to extract simplicity? Any examples?

Example: programming

Machine languages: no abstractions

- Hard to deal with low-level details
- Mastering complexity is crucial

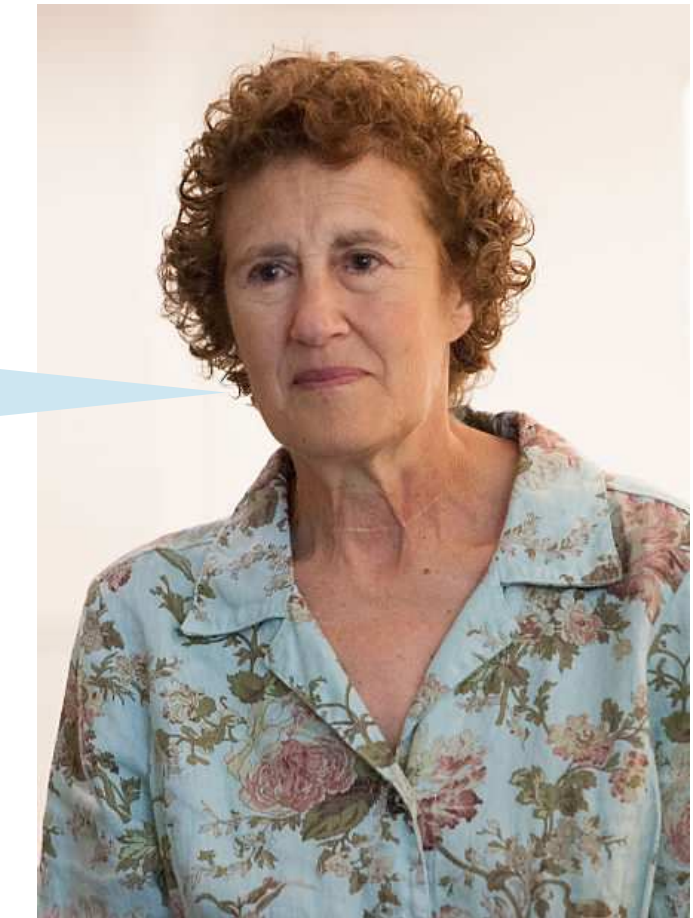
High-level languages: operating systems and other abstractions

- File systems, virtual memory, abstract data types...

Modern languages: even more abstractions

- Object oriented, garbage collection...

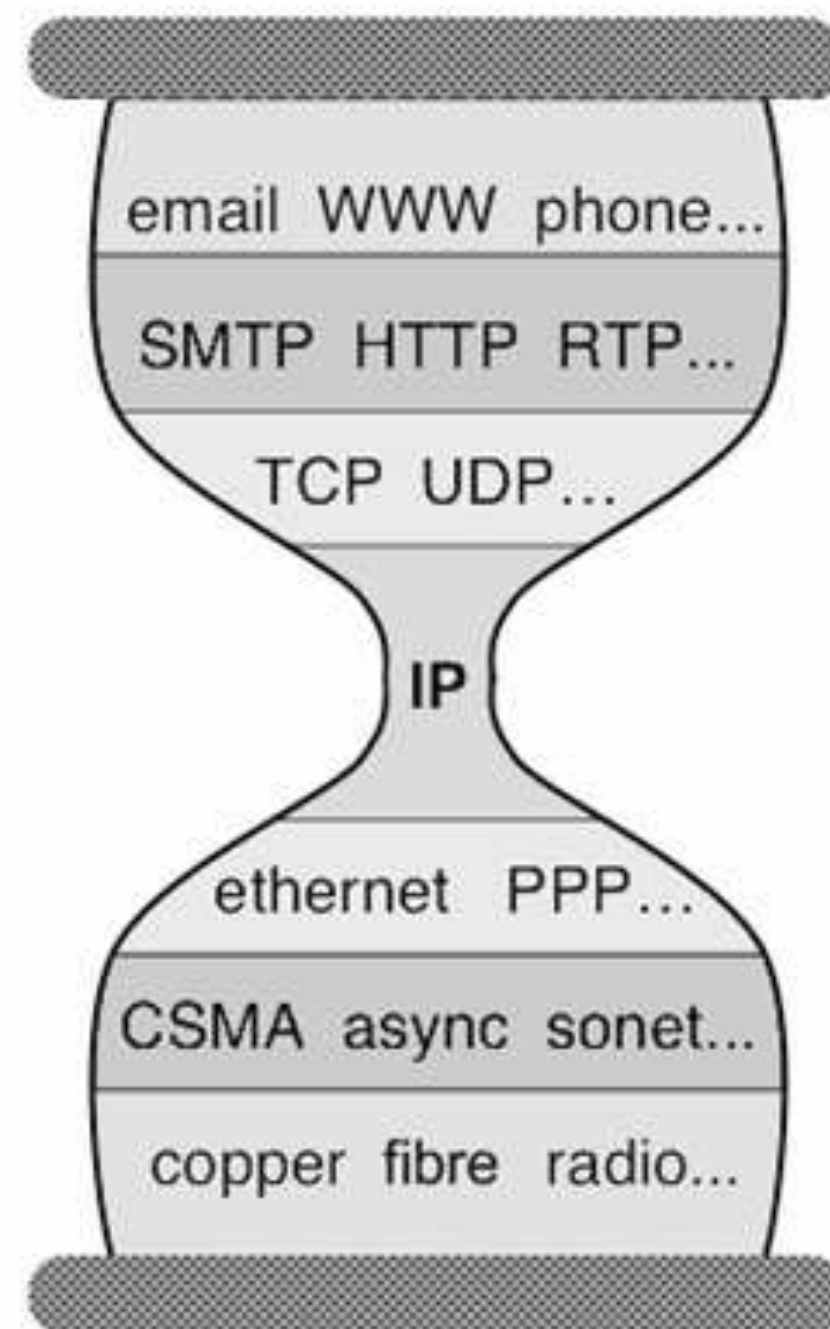
"Modularity based on abstractions is the way things get done!"



Barbara Liskov
(MIT, ACM Turing Award 2008,
pioneer in programming
languages, operating systems,
distributed computing)

What abstractions do we have in networking?

Abstractions for data plane



Applications, built on...

Reliable (or unreliable) transport, built on...

Best-effort global packet delivery, built on...

Best-effort local packet delivery, built on...

Physical transfer of bits

What about the control plane? Any abstractions?

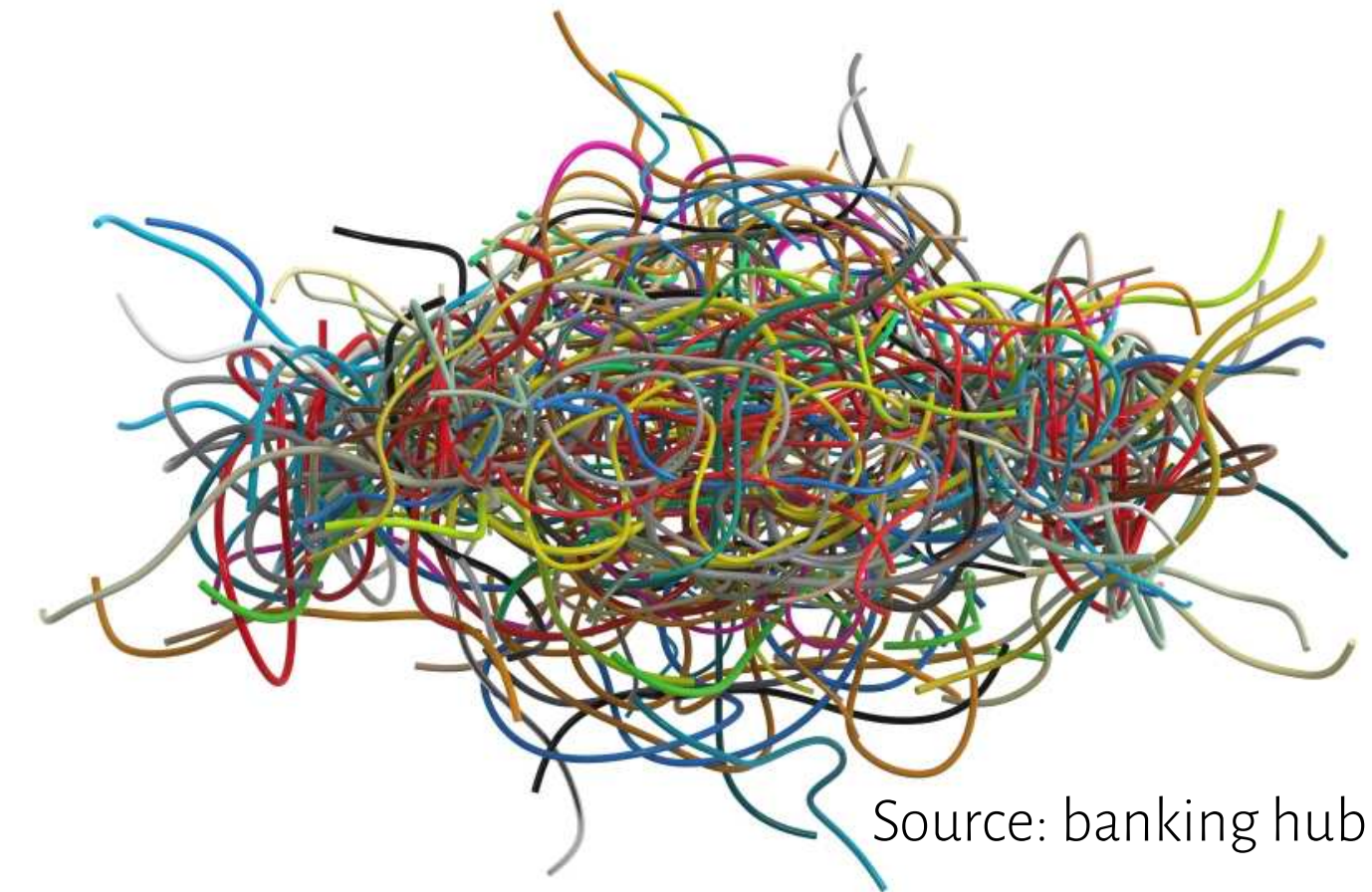
Control plane is full of mechanisms

Variety of goals, **no modularity**

- Routing: distributed routing algorithms
- Isolation: ACLs, VLANs, Firewalls,
- Traffic engineering: adjusting weights, MPLS,...

Control plane: **mechanism without abstraction**

You probably feel that computer networking is all about artifacts (e.g., artifacts, tools), and seems not a real discipline... (at least not like computer architecture, operating systems)



Source: banking hub

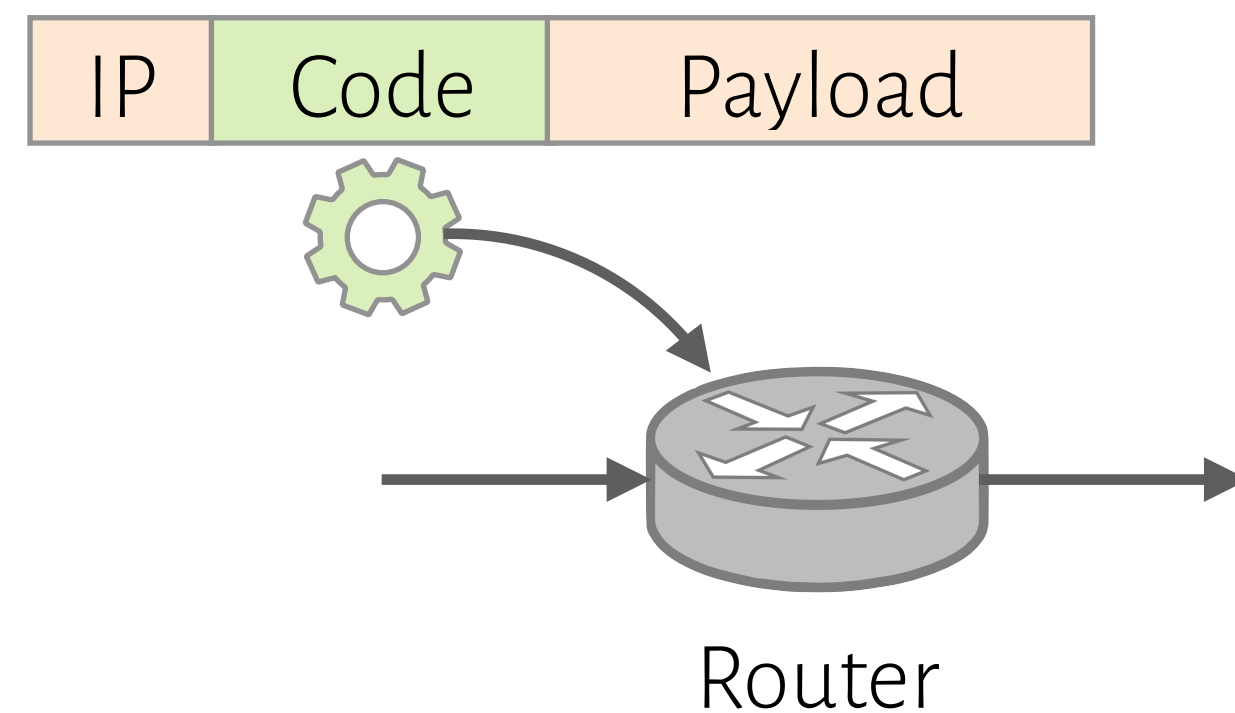
We need abstractions and ultimately, we should be able to **program the network** as we do for computers.

Questions?

The evolution: active networking (1990s)

First attempt making networks **programmable**: demultiplexing packets to software programs

Packet



In-band approach: The packet encapsulates a small piece of code that can be executed on the router, based on which the router decides what to do with the packet

Out-band approach: User injects the code to be executed beforehand → the programmable network approach which received a lot of attention recently.

We will discuss it in our next lecture!

The evolution: control/data plane separation (2003-2007)

4D (ACM SIGCOMM CCR 2004)

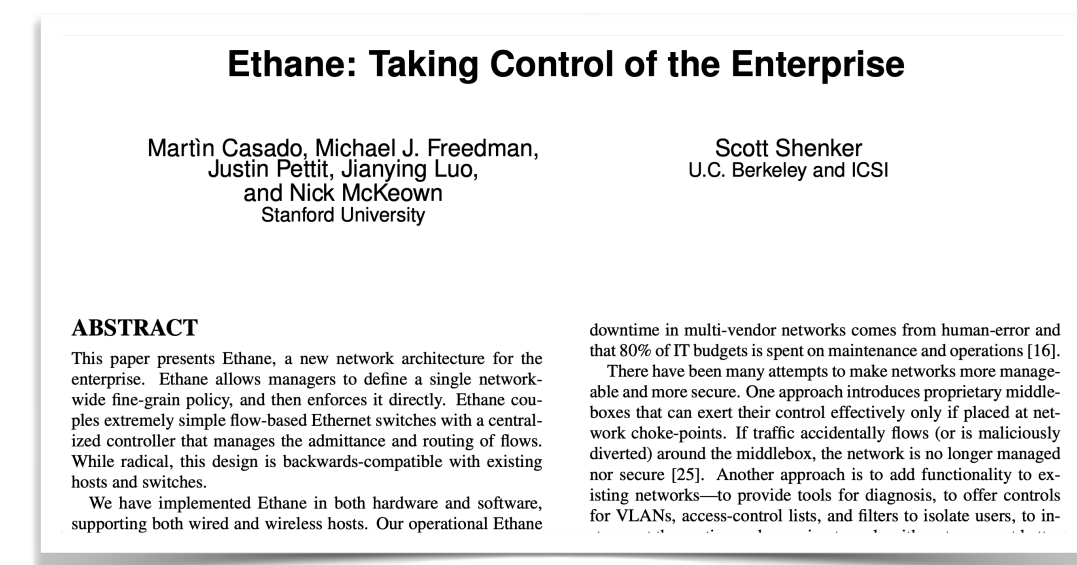
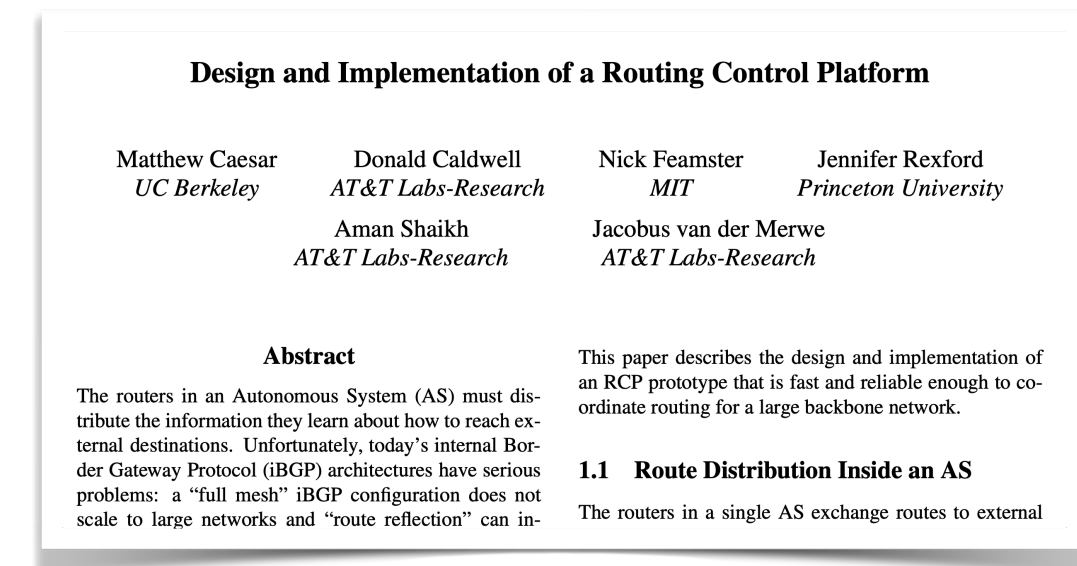
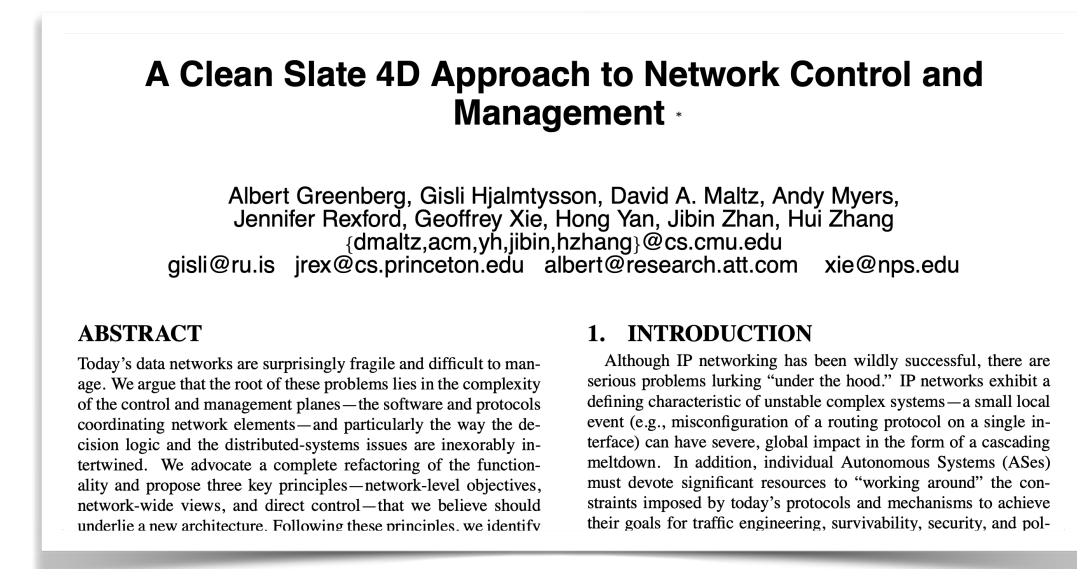
- Data, discovery, dissemination, decision
- Clean-slate: network-wide view, direct control, network-global objectives

RCP (USENIX NSDI 2005)

- Routing Control Platform for centralized inter-AS routing, replacing iBGP

Ethane (ACM SIGCOMM 2007)

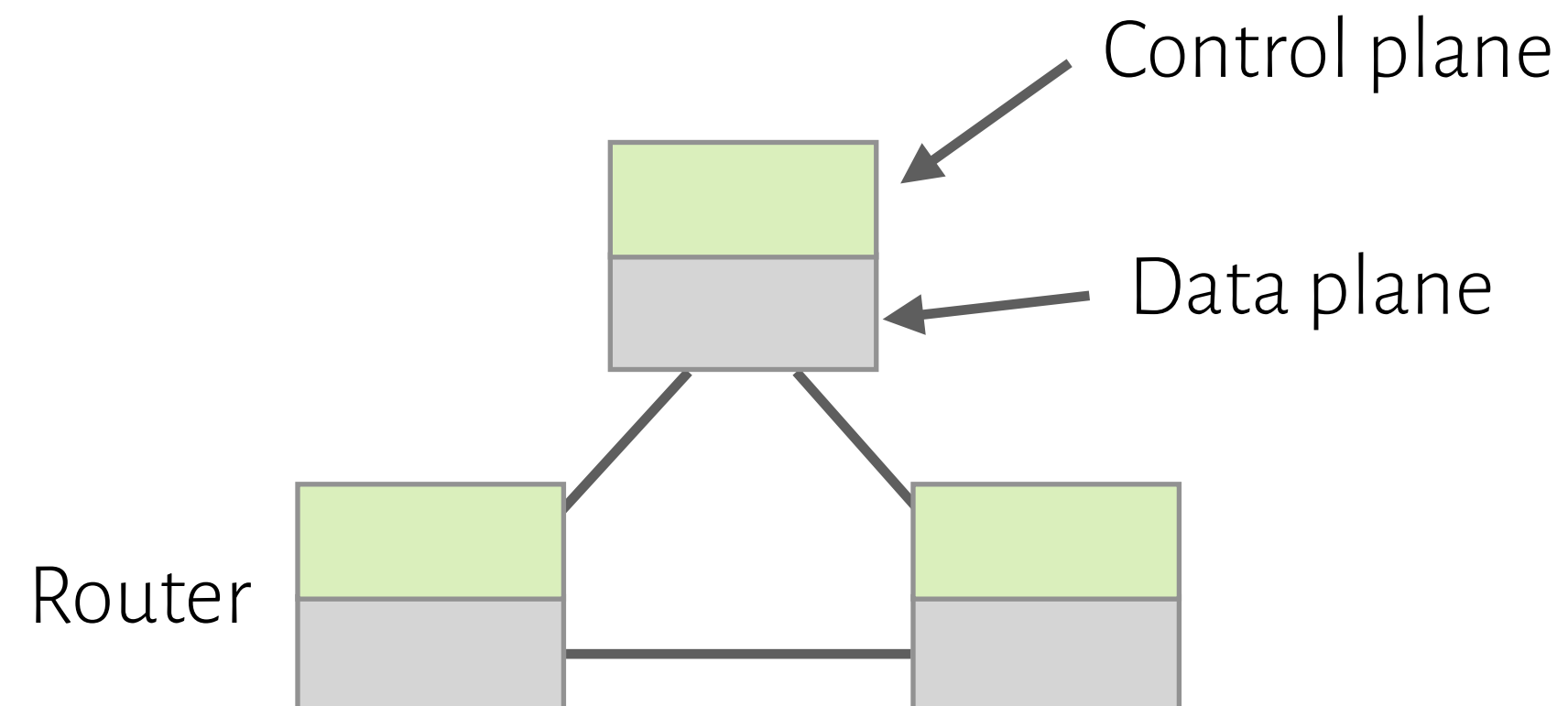
- Flow-based switching with centralized control for enterprise
- Precursor of SDN



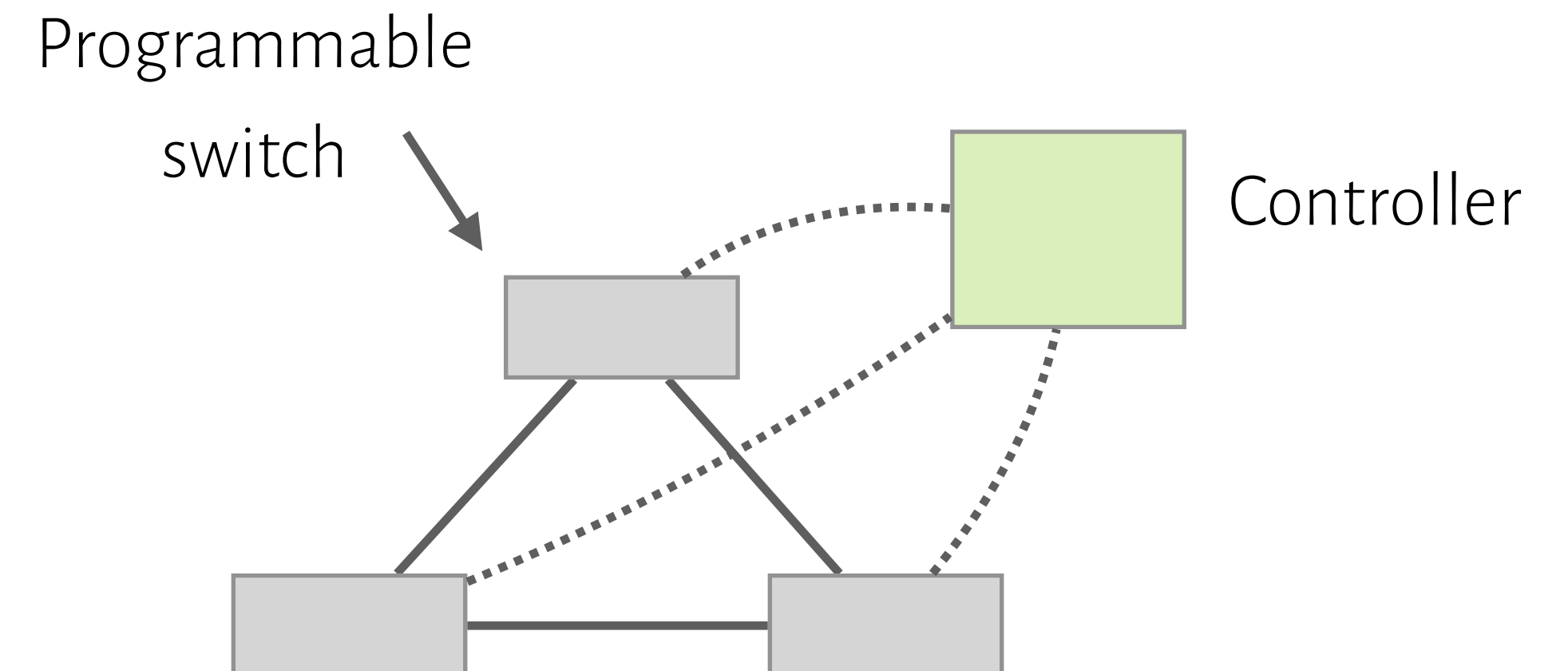
Software defined network

A network in which

- The **control** plane is physically **separate from** the **data** plane
- A **single** (logically centralized) **control** plane controls several forwarding devices



Traditional network



Software define network

The Road to SDN: An Intellectual History of Programmable Networks

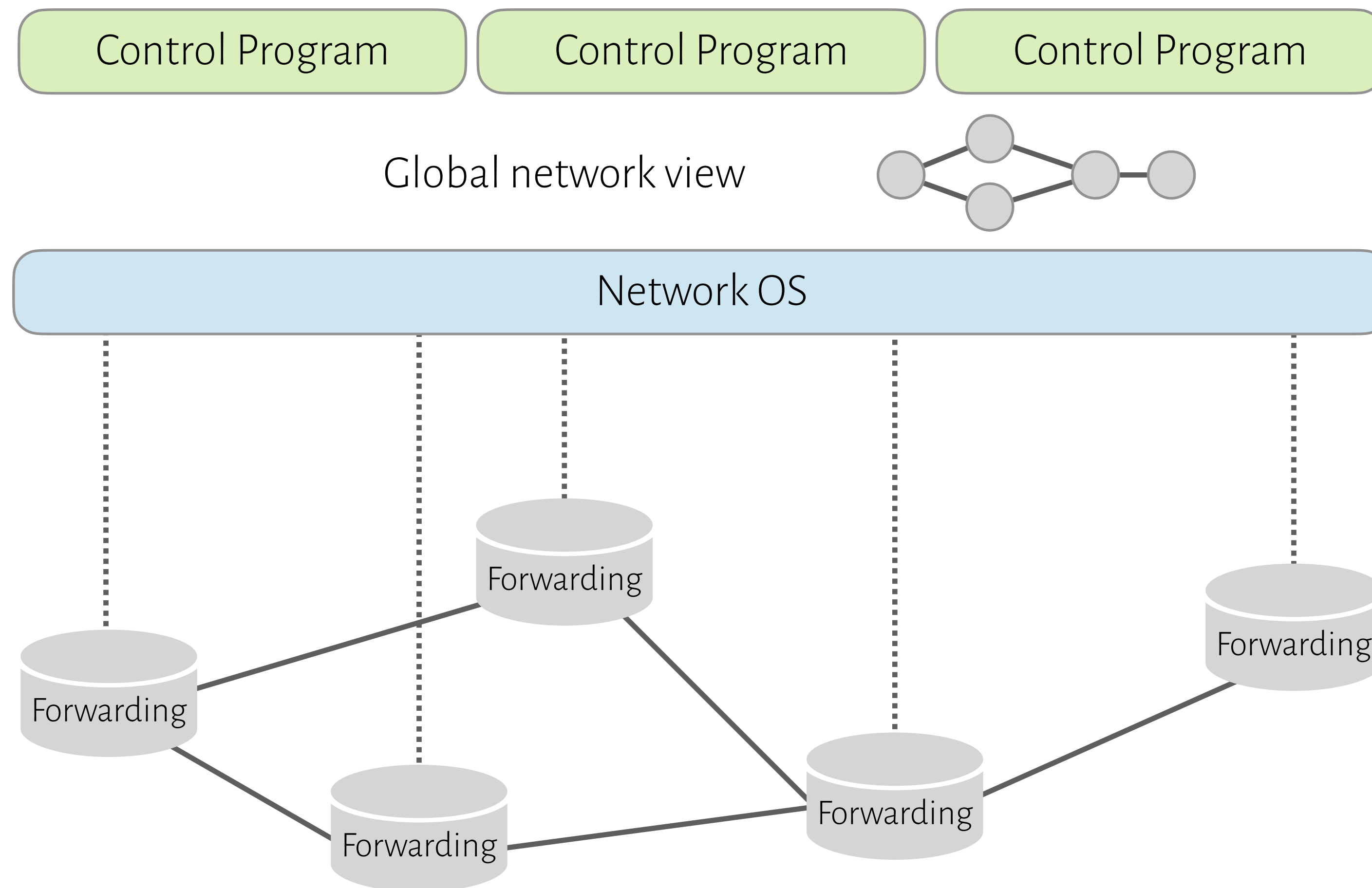
Nick Feamster Jennifer Rexford Ellen Zegura
Georgia Tech Princeton University Georgia Tech
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ABSTRACT

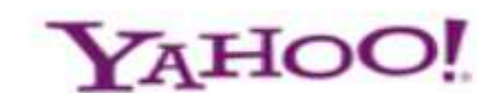
Software Defined Networking (SDN) is an exciting technology that enables innovation in how we design and manage networks. Although this technology seems to have appeared suddenly, SDN is part of a long history of efforts to make computer networks more programmable. In this paper, we trace the intellectual history of programmable networks, including active networks, early efforts to separate the control and data plane, and more recent work on OpenFlow and network operating systems. We highlight key concepts, as well as the technology pushes and application pulls that spurred each innovation. Along the way, we debunk common myths and misconceptions about the technologies and clarify the relationship between SDN and related technologies such as network virtualization.

makes). Second, an SDN consolidates the control plane, so that a single software control program controls *multiple* data-plane elements. The SDN control plane exercises direct control over the state in the network's data-plane elements (*i.e.*, routers, switches, and other middleboxes) via a well-defined Application Programming Interface (API). OpenFlow [51] is a prominent example of such an API. An OpenFlow switch has one or more tables of packet-handling rules. Each rule matches a subset of traffic and performs certain actions on the traffic that matches a rule; actions include dropping, forwarding, or flooding. Depending on the rules installed by a controller application, an OpenFlow switch can behave like a router, switch, firewall, network address translator, or something in between.

SDN architecture overview

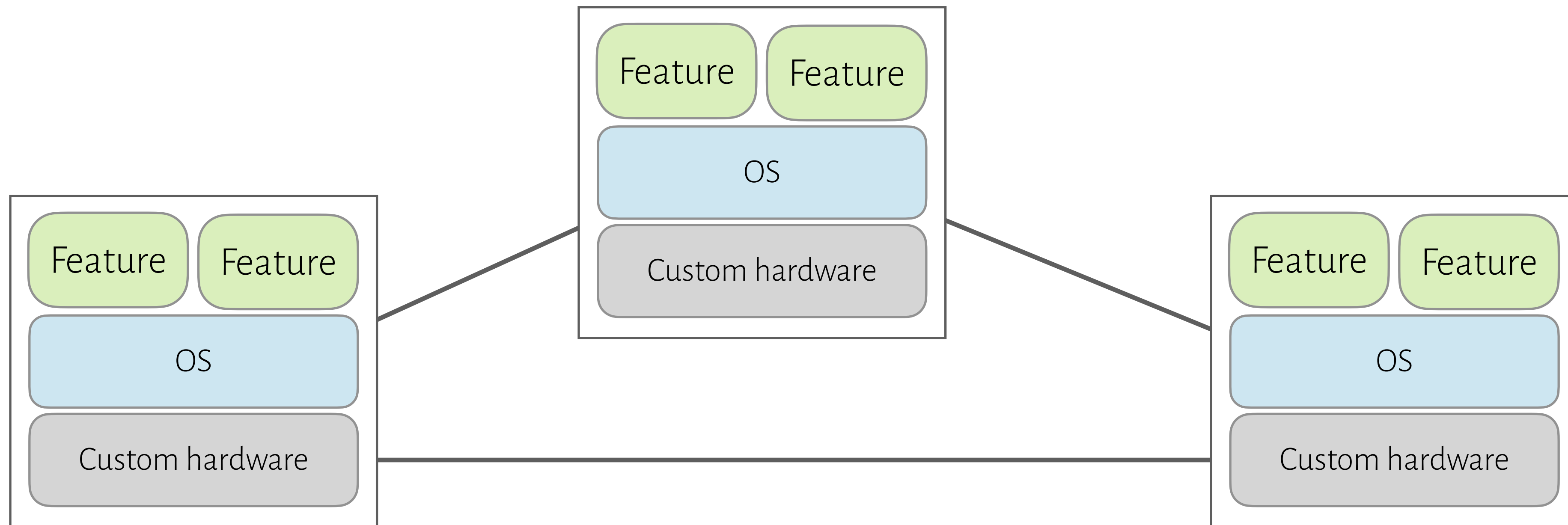


A major trend in networking

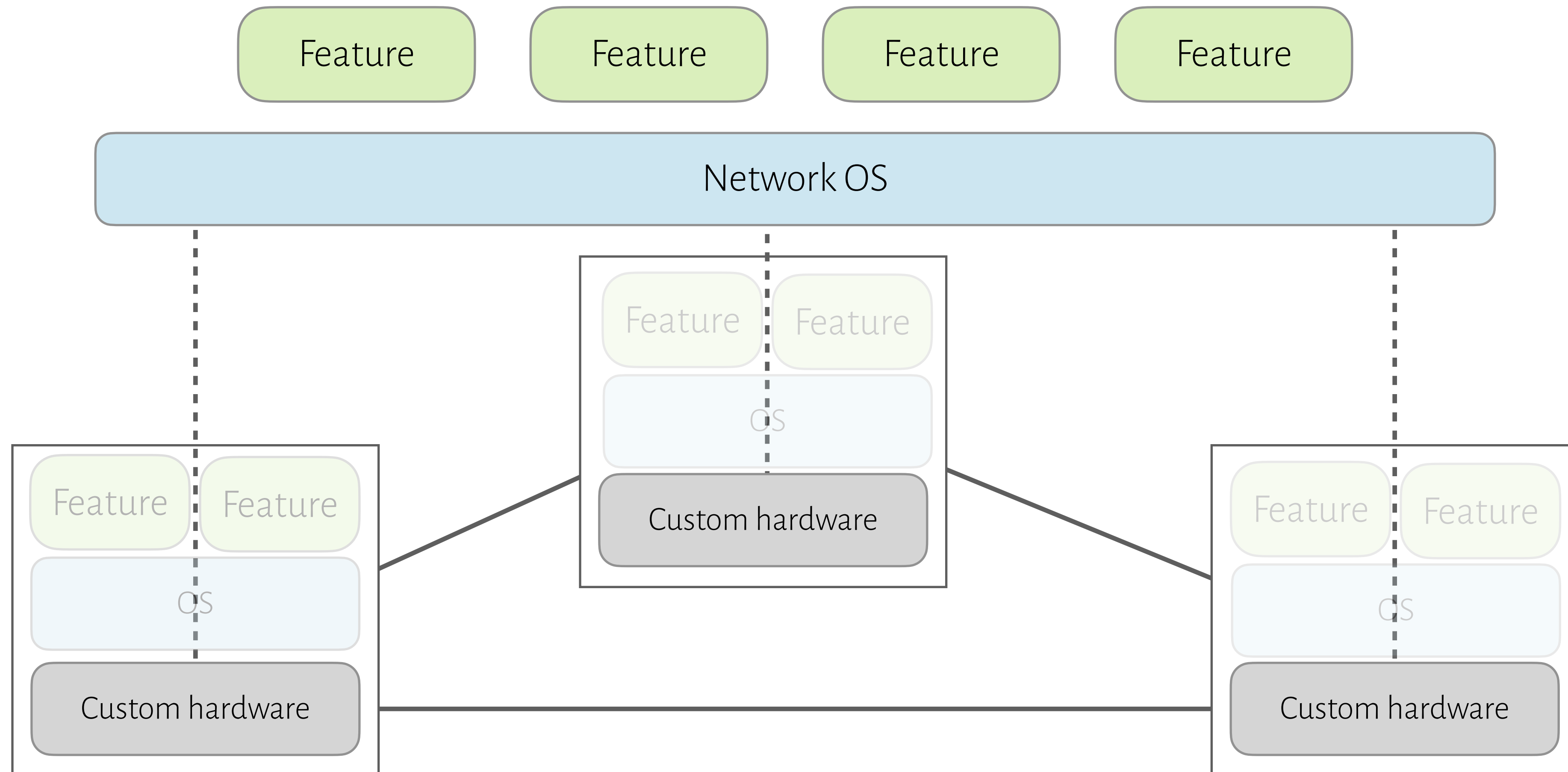


Startup from Stanford and Berkeley in 2007, acquired by VMware in 2012 for \$1.26 billion.

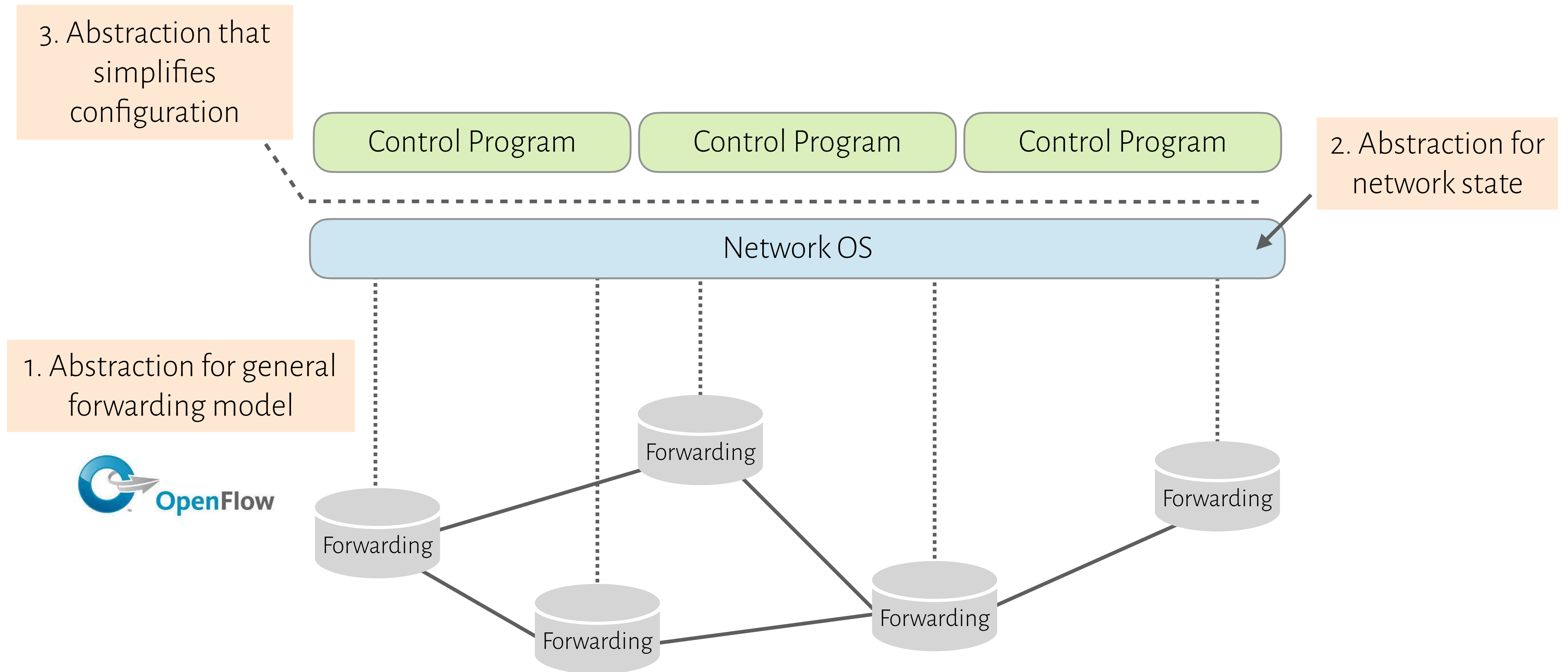
How SDN changes the network



How SDN changes the network



Abstractions in SDN



Abstraction #1: forwarding abstraction

Express intent independent of implementation

OpenFlow is the current proposal for forwarding

- Standardized interface to switch: non-proprietary COTS hardware and software
- Configuration in terms of flow entries: <header, action>
- No hardware modifications needed, simply a firmware update

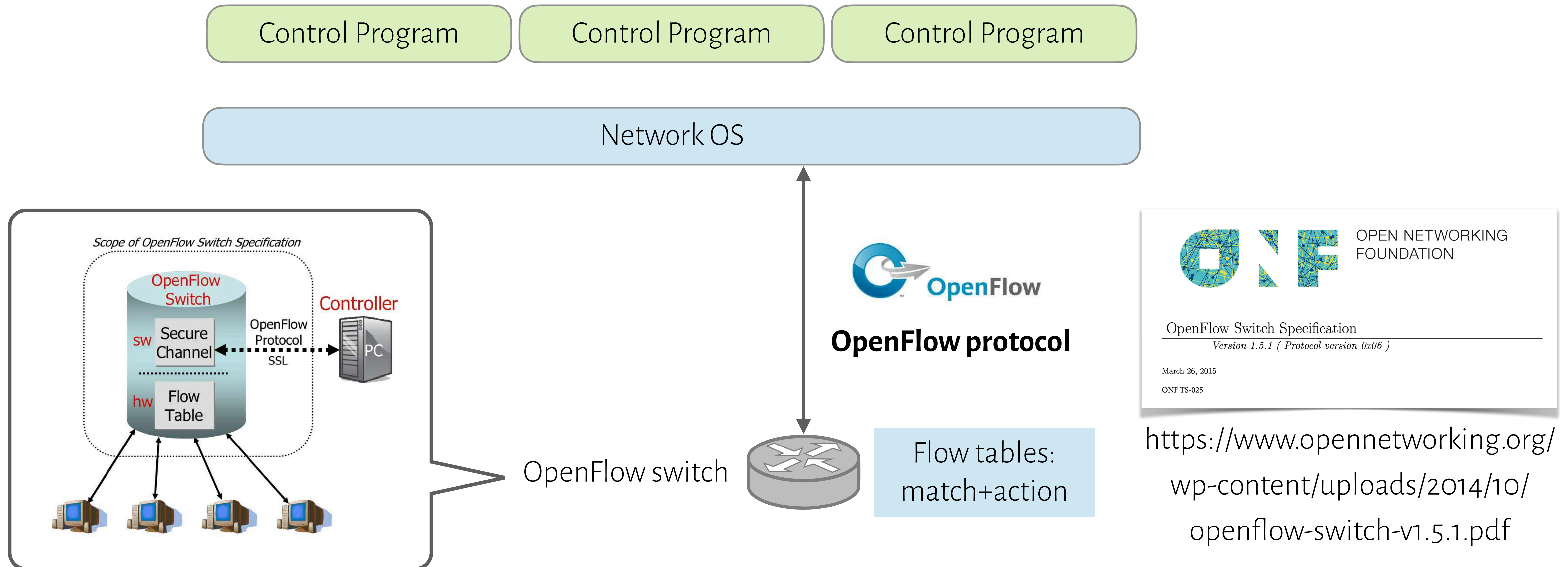
Design details concern exact nature of **match+action**

Benefits

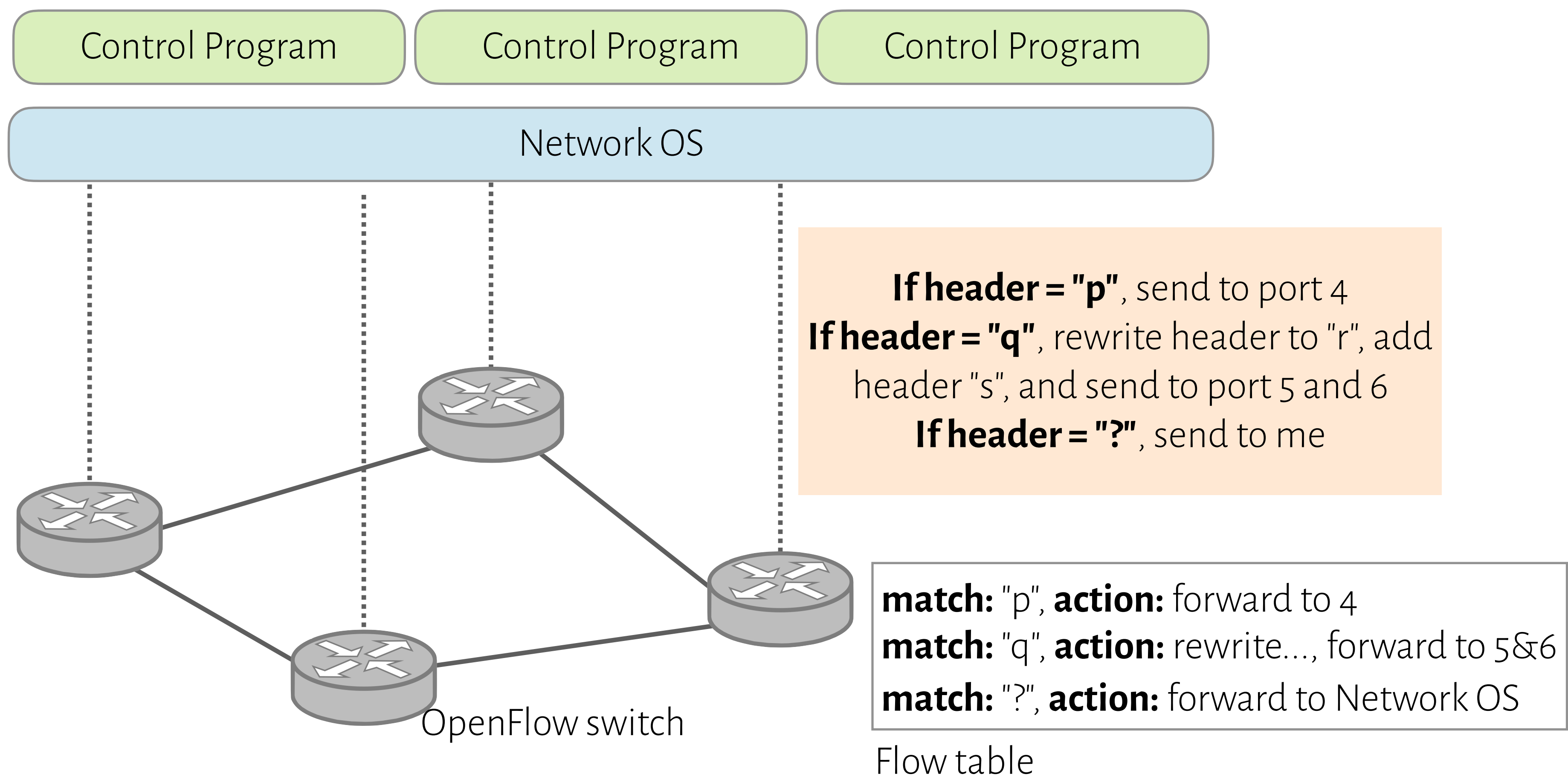
- Much cheaper, no more \$27K for a single switch
- No vendor lock-in



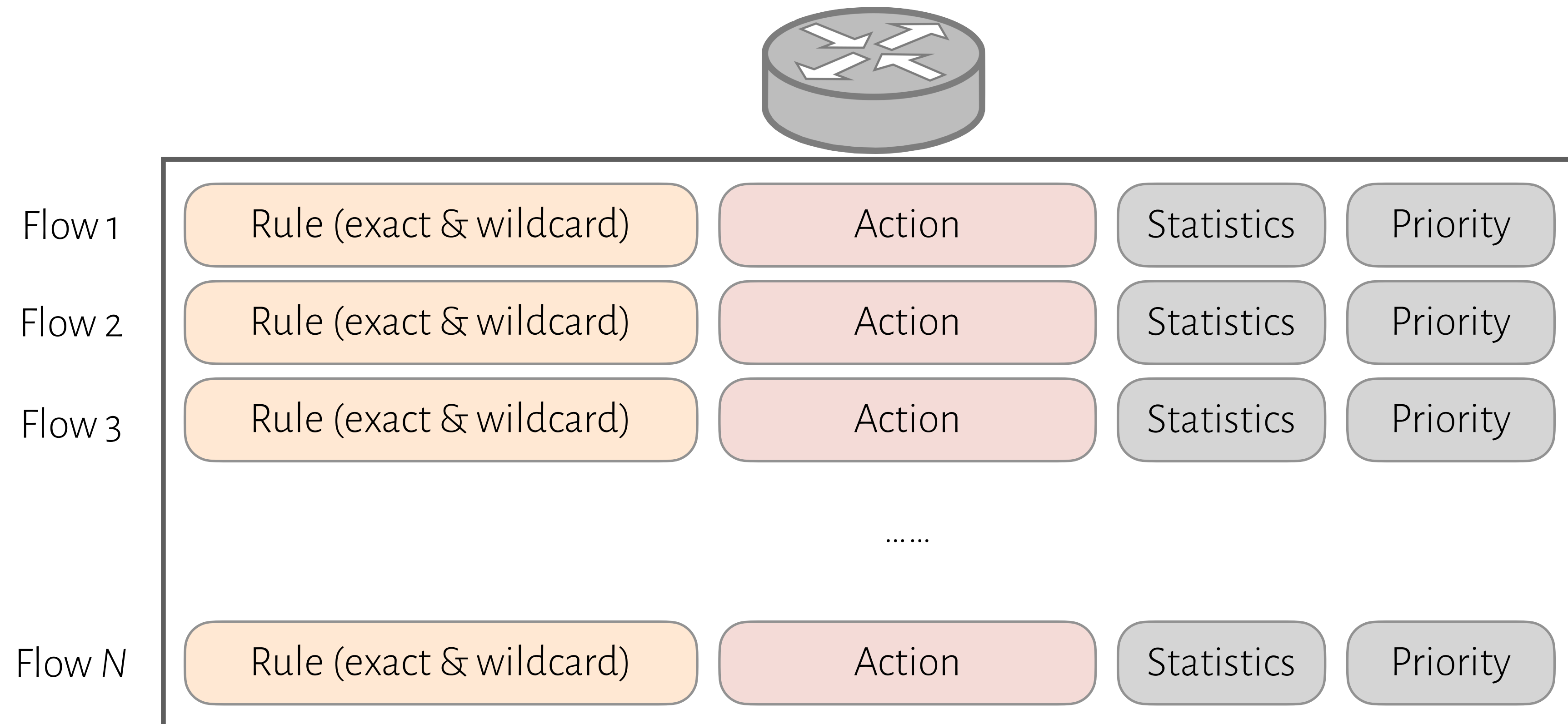
OpenFlow



OpenFlow example



Flow table(s) on OpenFlow switches



Exploit the forwarding tables that are already in routers, switches, and chipsets

Match+action

Match arbitrary bits in headers

- Match on any header, or new header
- Allows any flow granularity

Action

- Forward to port(s), drop, send to the controller
- Overwrite header with mask, push or pop
- Forward at specific bit-rate
- Do not support payload-related network functions like deep packet inspection

In Port	VLAN ID	Ethernet			IP			TCP	
		SA	DA	Type	SA	DA	Proto	Src	Dst



Match: 1000X01XX0101001X

Abstraction #2: network state abstraction

Global network view

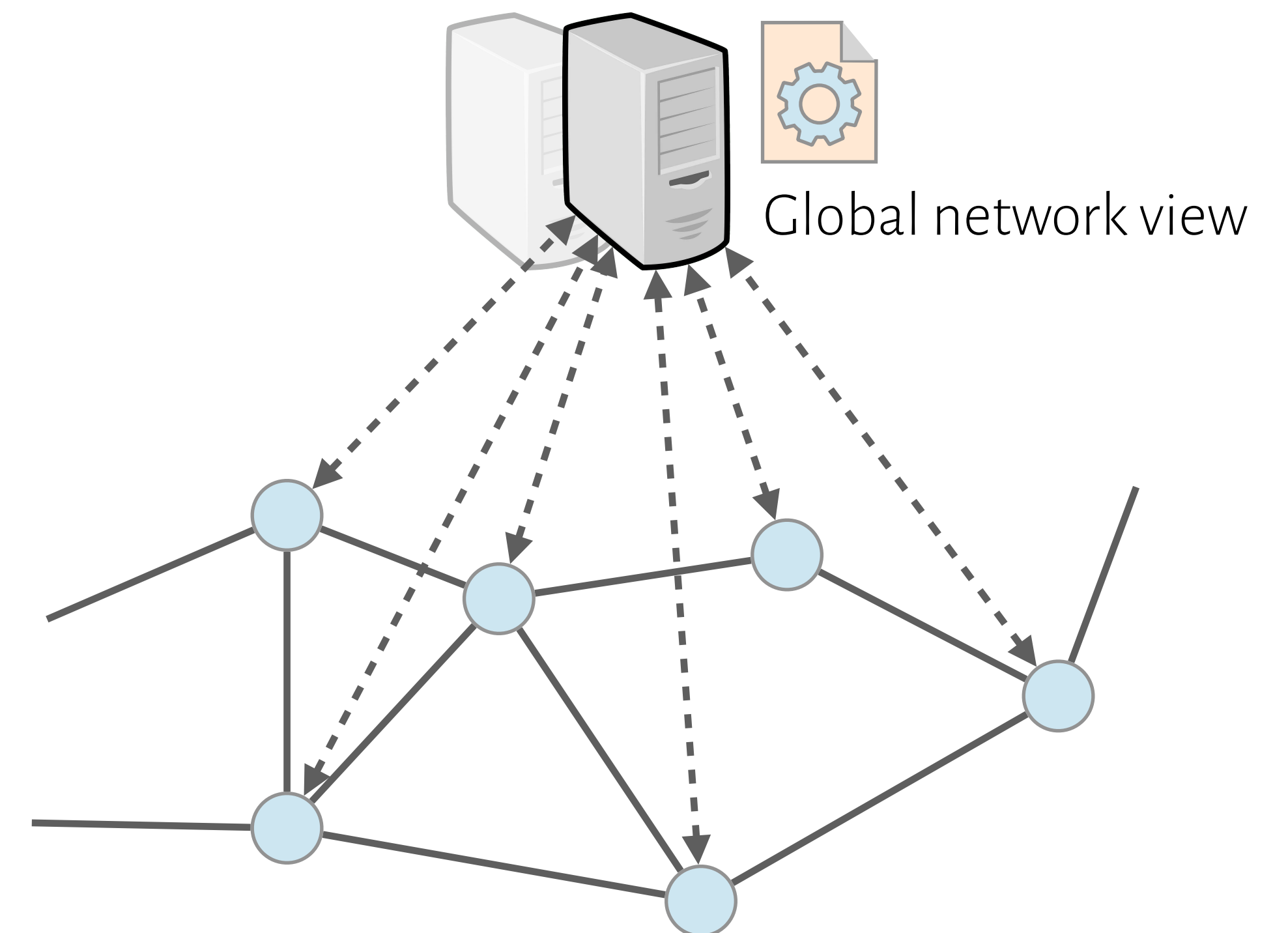
- Annotated network graph provided through an API

Implementation: "Network Operating Systems"

- Runs on servers in network (as "controllers")
- Replicated for reliability

Information flows both ways

- Information from routers/switches to form view
- Configurations to routers/switches to control forwarding

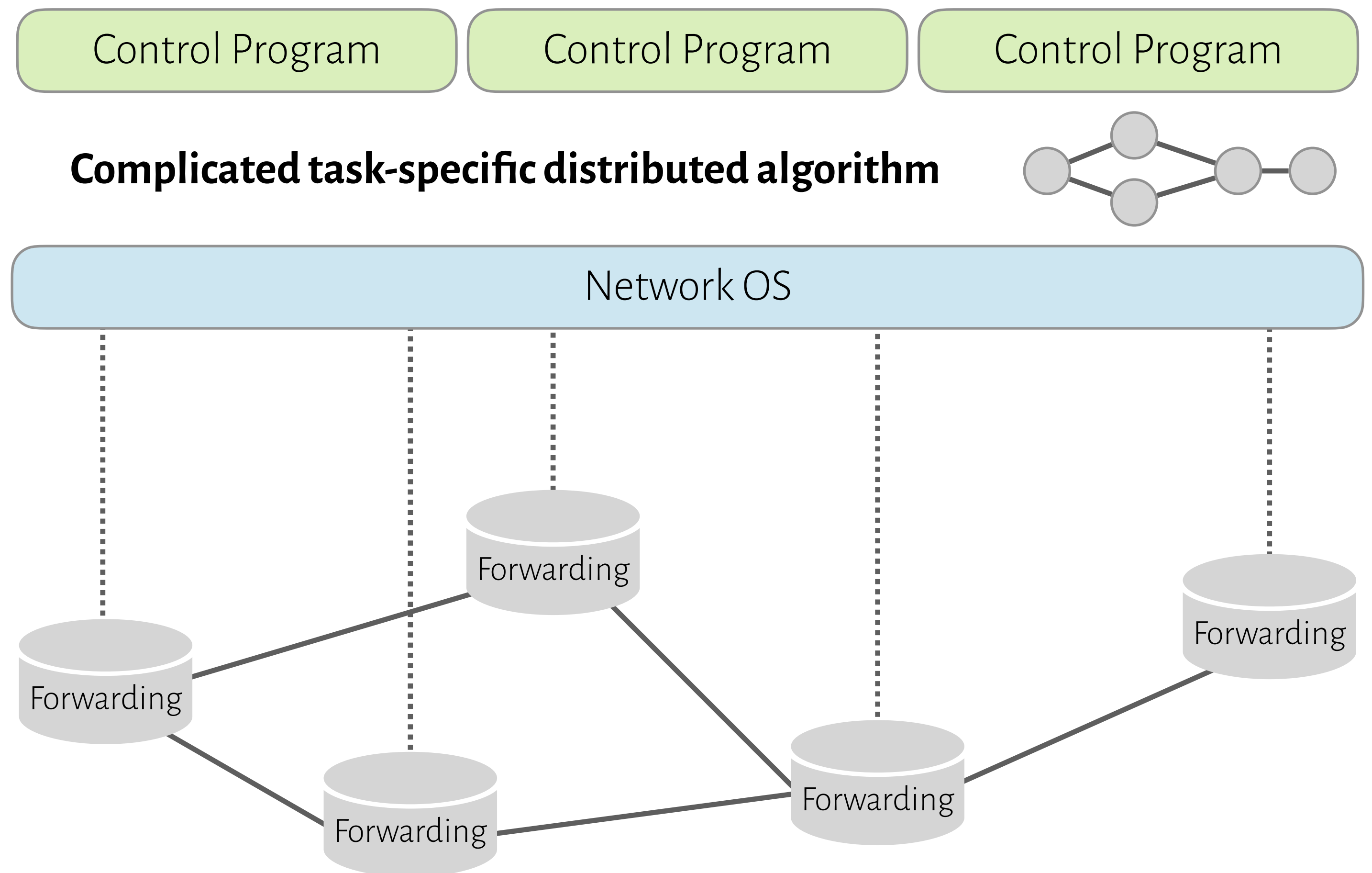


Major change in paradigm

Control program: Configuration =
Function(View)

Control mechanism now is a
program using Network OS APIs

Not a distributed protocol, just a
graph algorithm



Abstraction #3: specification abstraction

Control mechanism expresses **desired behavior**

- Whether it be isolation, access control, or QoS

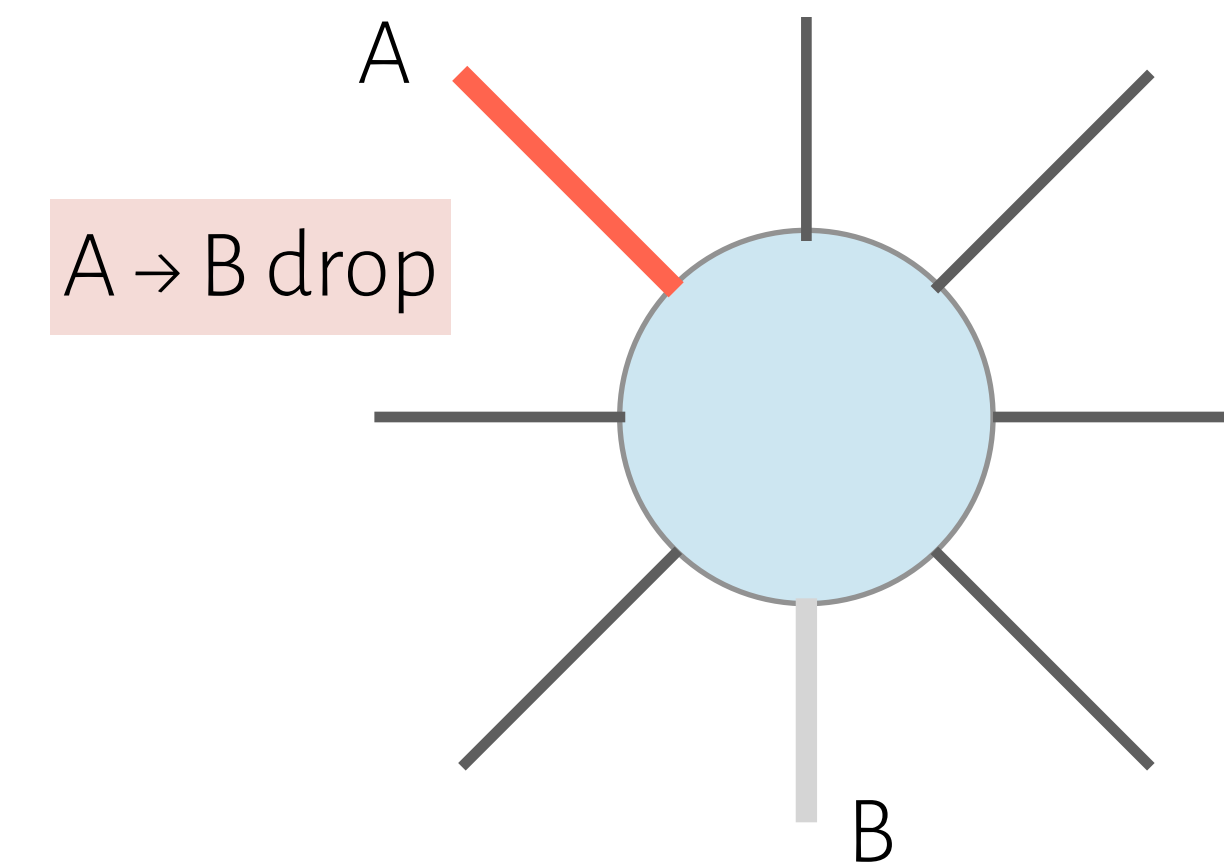
It should not be responsible for implementing that behavior on physical network infrastructure

- Requires configuring the forwarding tables in each switch

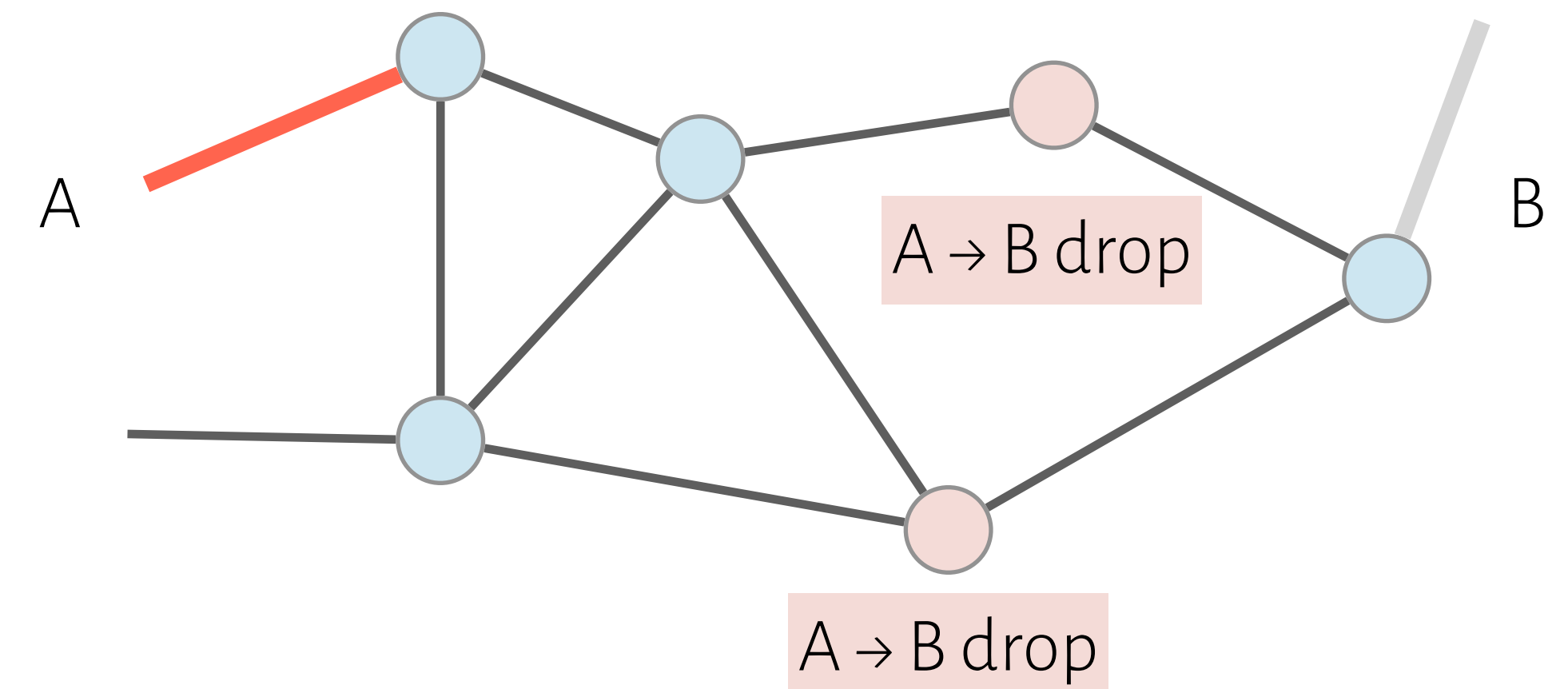
Proposed abstraction: **abstract view of the network**

- Abstract view models only enough detail to specify goals
- Will depend on task semantics

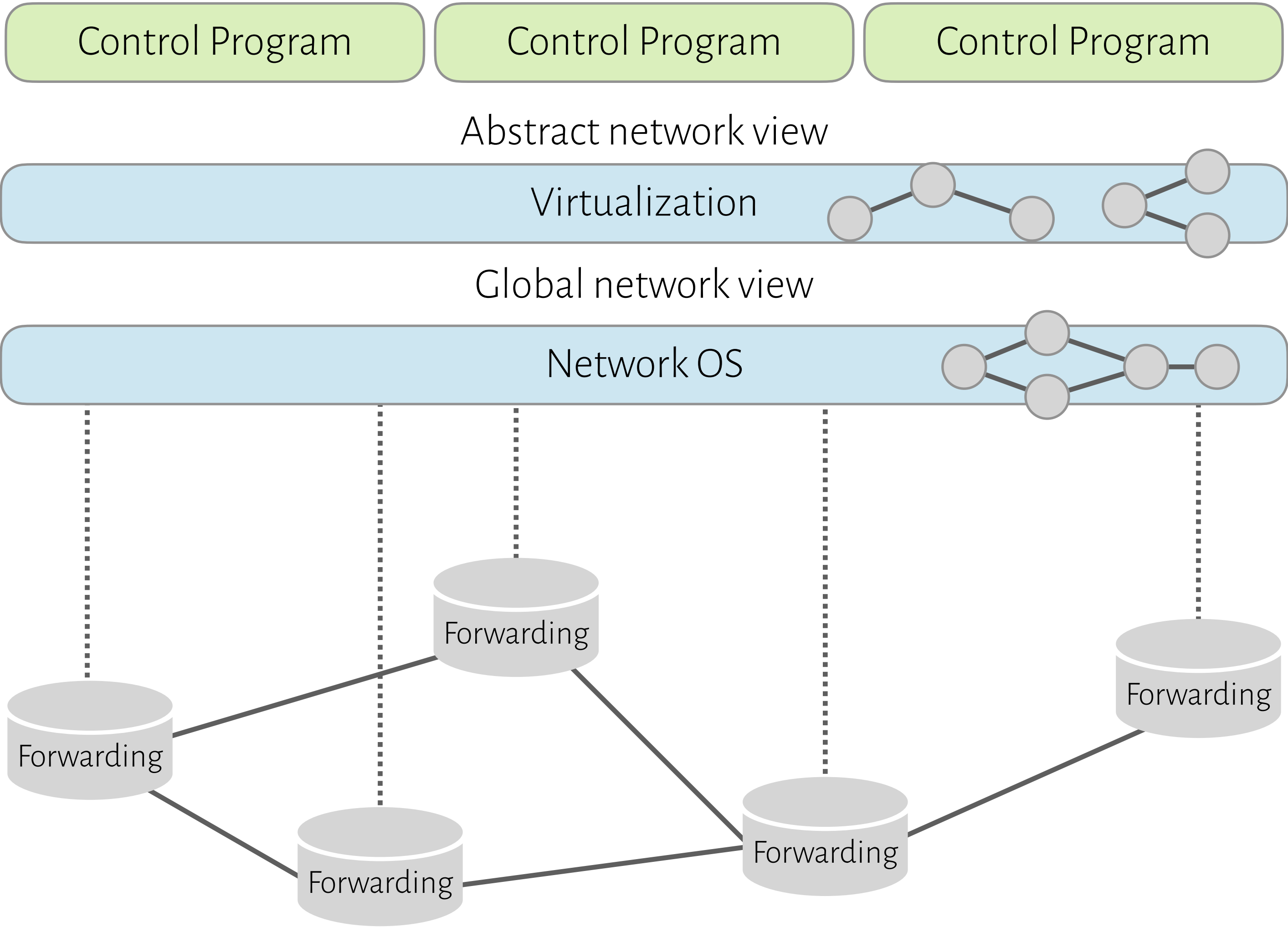
Abstract network view



Global network view



SDN control plane layers

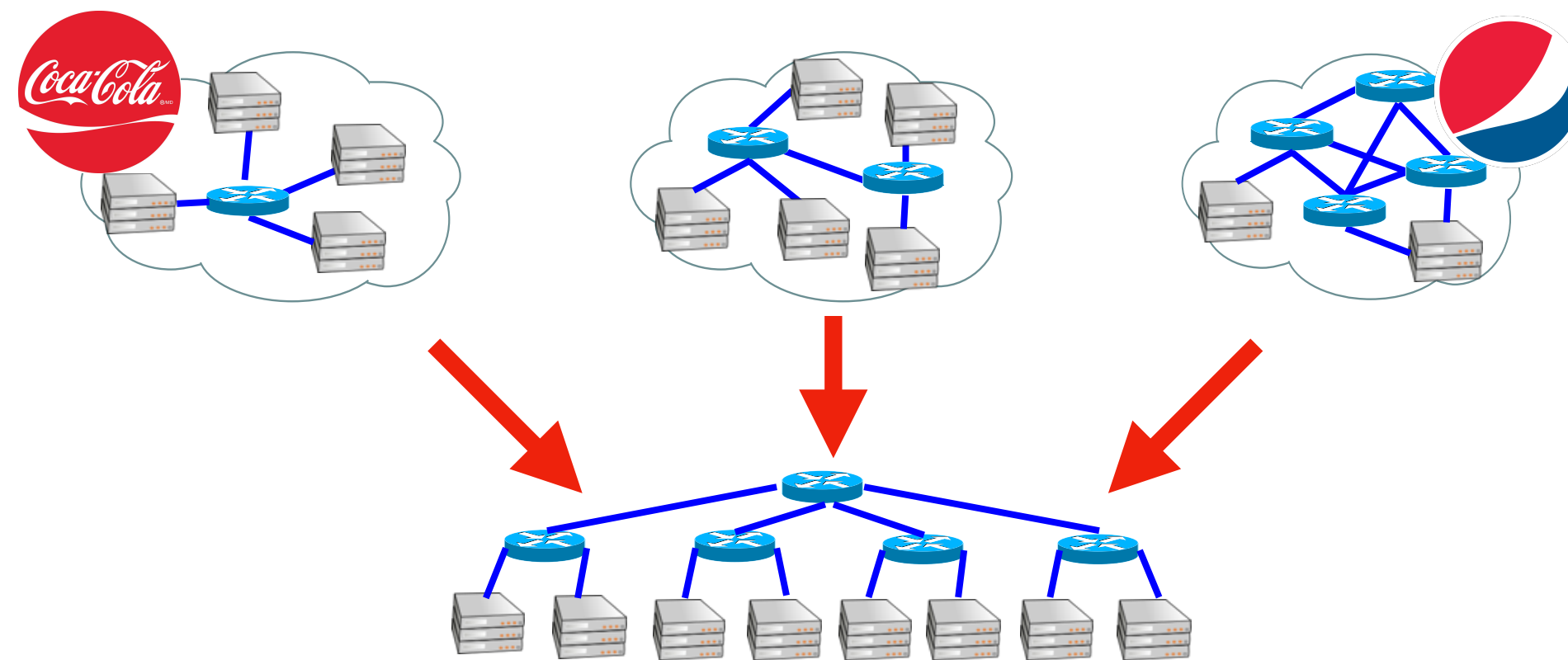


Questions?

Network virtualization

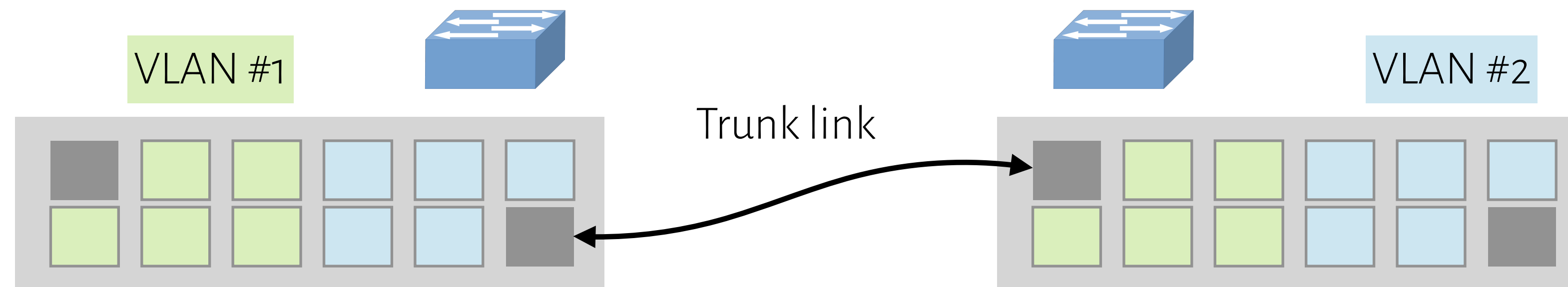
Ability to run multiple virtual networks that

- Each has a separate control and data plane
- Co-exist together on top of one physical network
- Can be managed by individual parties that potentially do not trust each other



What techniques we have learned that can be used for network virtualization?

VLAN



VLAN provides basic isolation for Ethernet LANs, but it has many problems for network virtualization

- **Cannot program packet forwarding:** stuck with learning switch and spanning tree
- **No obvious opt-in mechanisms:** who maps a packet to a VLAN?
- **Resource isolation** is problematic

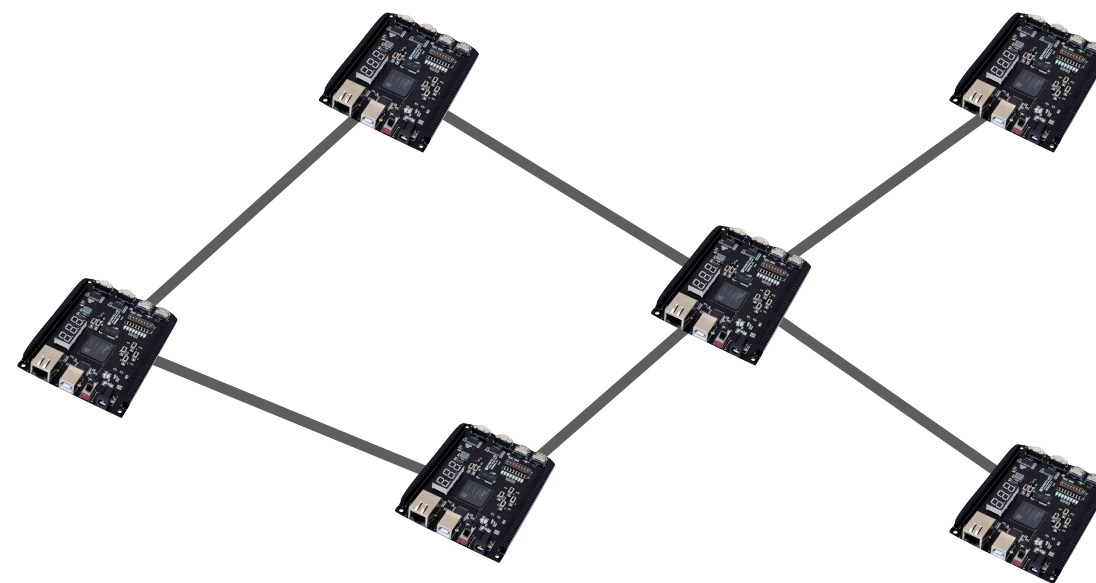
Network virtualization: use case

Imagine you come up with a novel network service, e.g., a new routing protocol, network load-balancer, how would you convince people that this is useful?

Network virtualization: use case

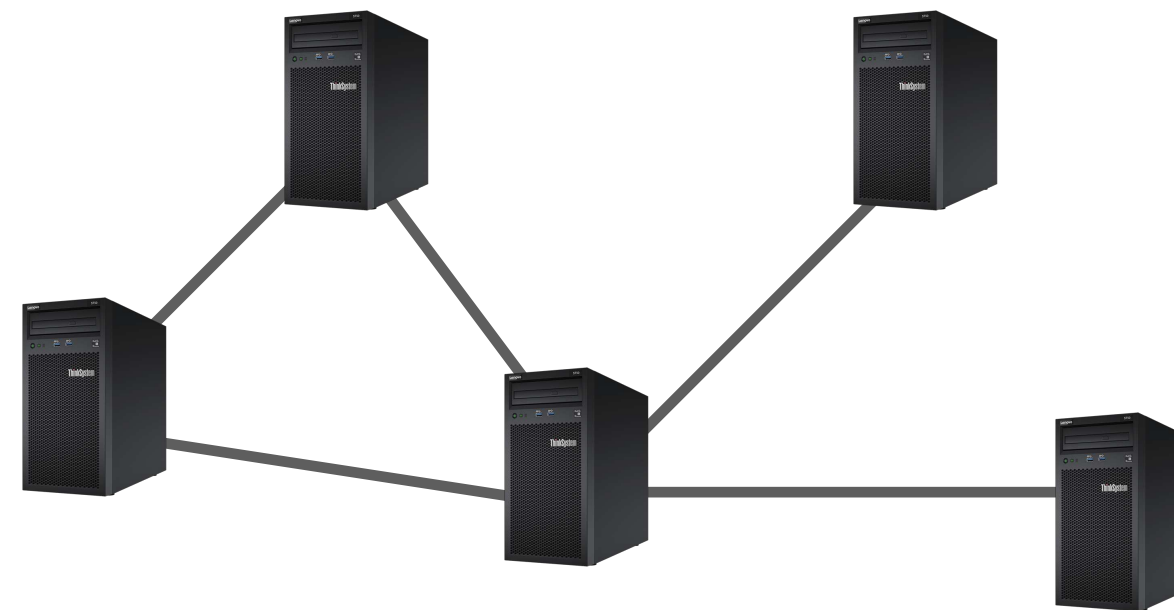
Imagine you come up with a novel network service, e.g., a new routing protocol, network load-balancer, how would you convince people that this is useful?

Hardware testbed



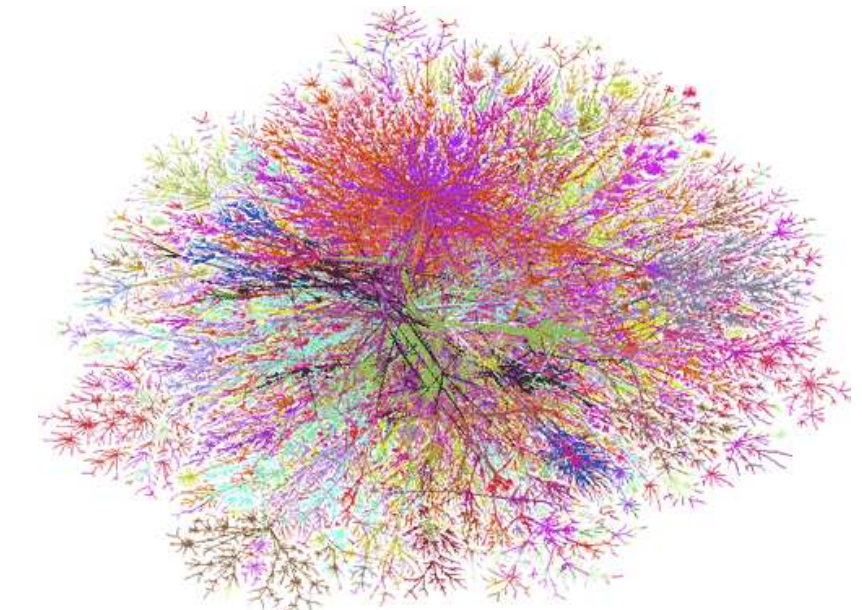
Expensive! Small-scale (fanout is small due to limited port number on NetFPGA)!

Software testbed



Large-scale (VINI/PlanetLab, Emulab)
Performance is slow (CPU-based), no realistic topology, hard to maintain!

Wild test on the Internet



Convincing network operators to try something new is very difficult! (Outages are the worst)

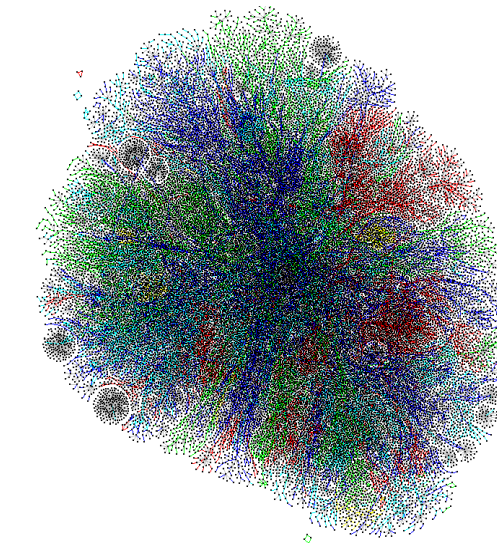
Problem

Realistically evaluating new network services is **hard**

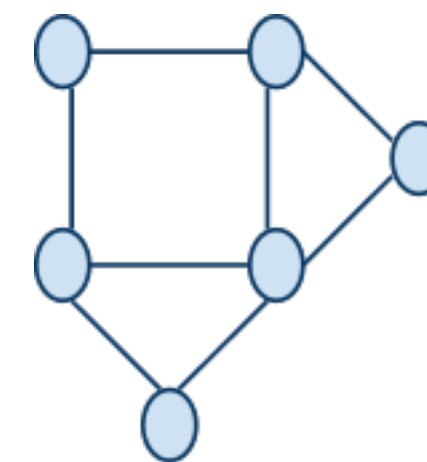
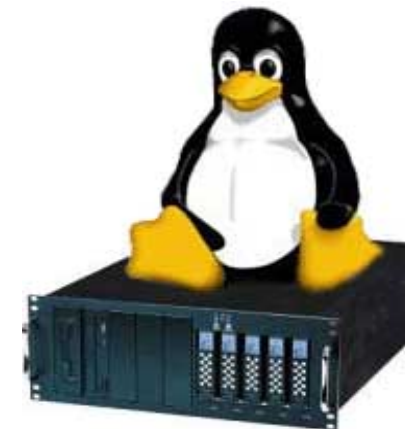
- Services that require changes to switches and routers
- For example: routing protocols, traffic monitoring services, IP mobility

Results

- Many good ideas do not get deployed
- Many deployed services still have bugs



Real networks



Test environments

Solution: network slicing

Divide the production network into logical **slices**

- Each slice/service controls its own packet forwarding
- Users pick which slice controls their traffic: opt-in
- Existing production services run in their own slice: spanning tree, OSPF/BGP

Enforce **strong isolation** between slices

- Actions in one slice do not affect others

Allow the (logical) **testbed** to mirror the **production** network

- Real hardware, performance, topologies, scale, users

Can the Production Network Be the Testbed?

Rob Sherwood*, Glen Gibb†, Kok-Kiong Yap†, Guido Appenzeller‡,
Martin Casado◇, Nick McKeown†, Guru Parulkar†

* Deutsche Telekom Inc. R&D Lab, Los Altos, CA† Stanford University, Palo Alto, CA

◇ Nicira Networks, Palo Alto, CA

‡ Big Switch Networks, Palo Alto, CA

Abstract

A persistent problem in computer network research is validation. When deciding how to evaluate a new feature or bug fix, a researcher or operator must trade-off realism (in terms of scale, actual user traffic, real equipment) and cost (larger scale costs more money, real user traffic likely requires downtime, and real equipment requires vendor adoption which can take years). Building a realistic testbed is hard because “real” networking takes place on closed, commercial switches and routers with special purpose hardware. But if we build our testbed from software switches, they run several orders of magnitude slower. Even if we build a realistic network testbed, it

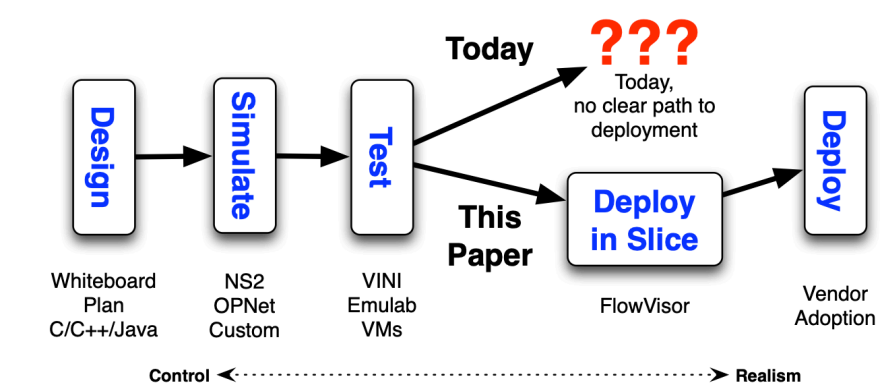
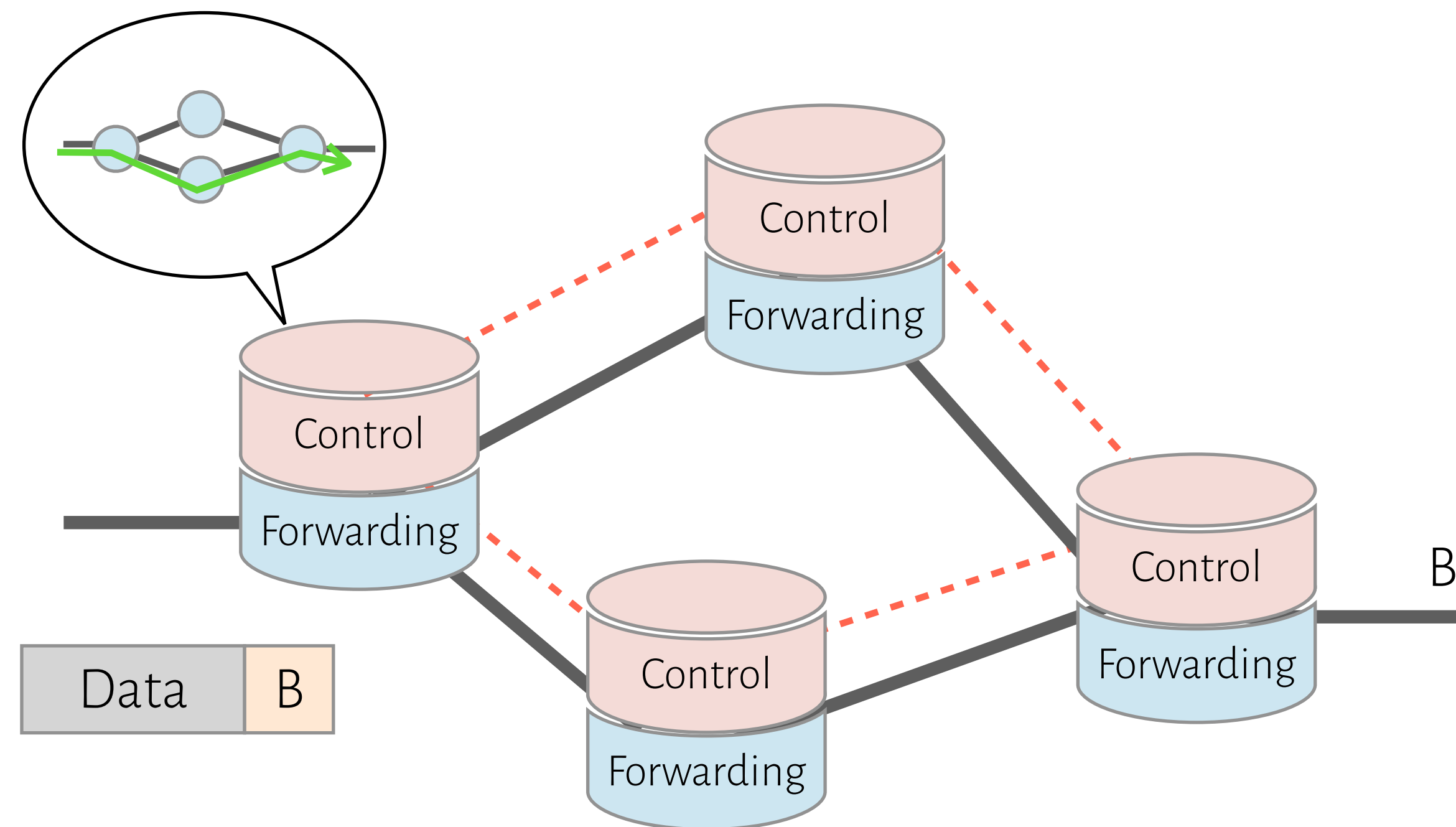


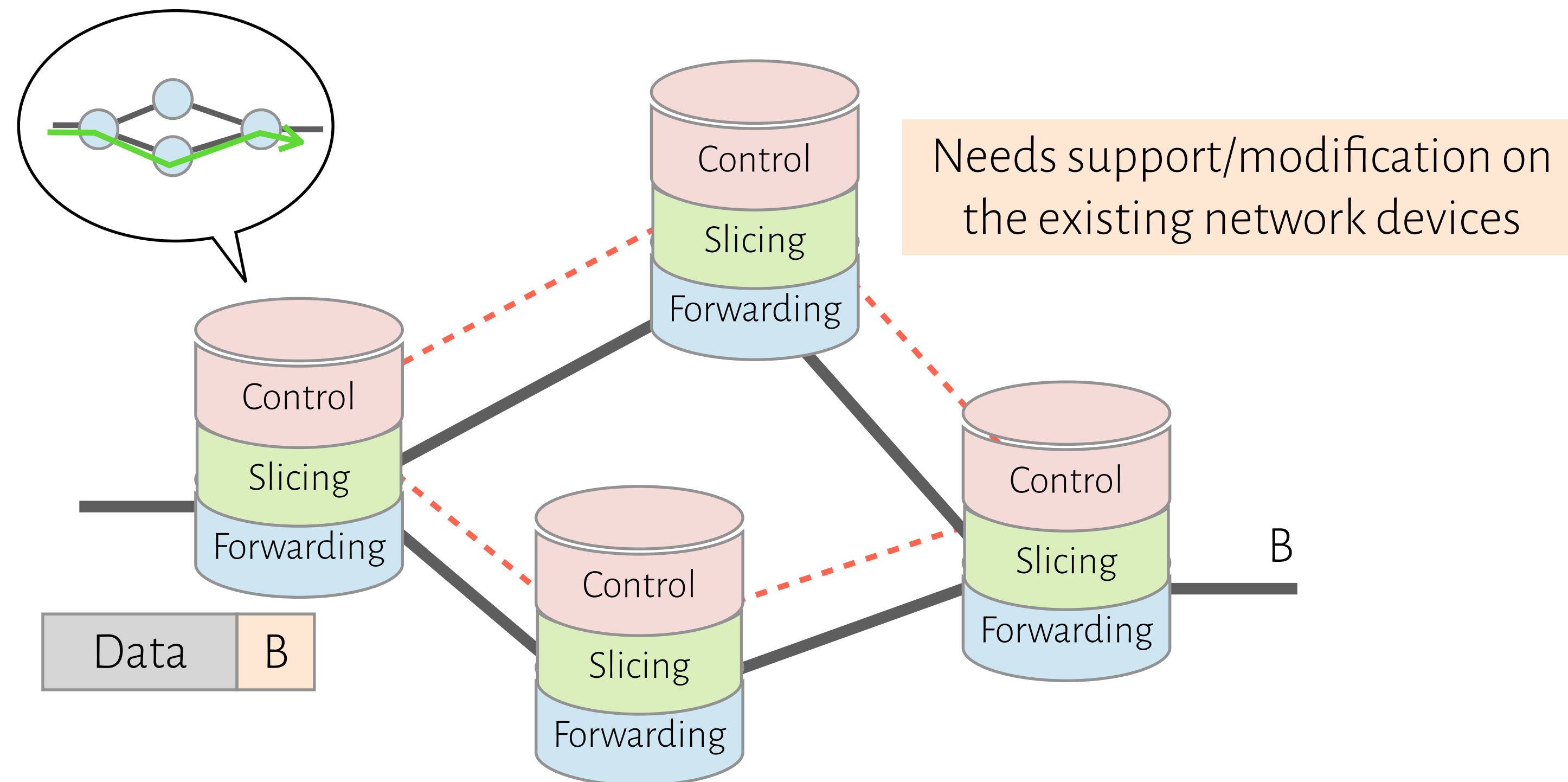
Figure 1: Today's evaluation process is a continuum from controlled but synthetic to uncontrolled but realistic testing, with no clear path to vendor adoption.

USENIX OSDI 2010

Traditional network

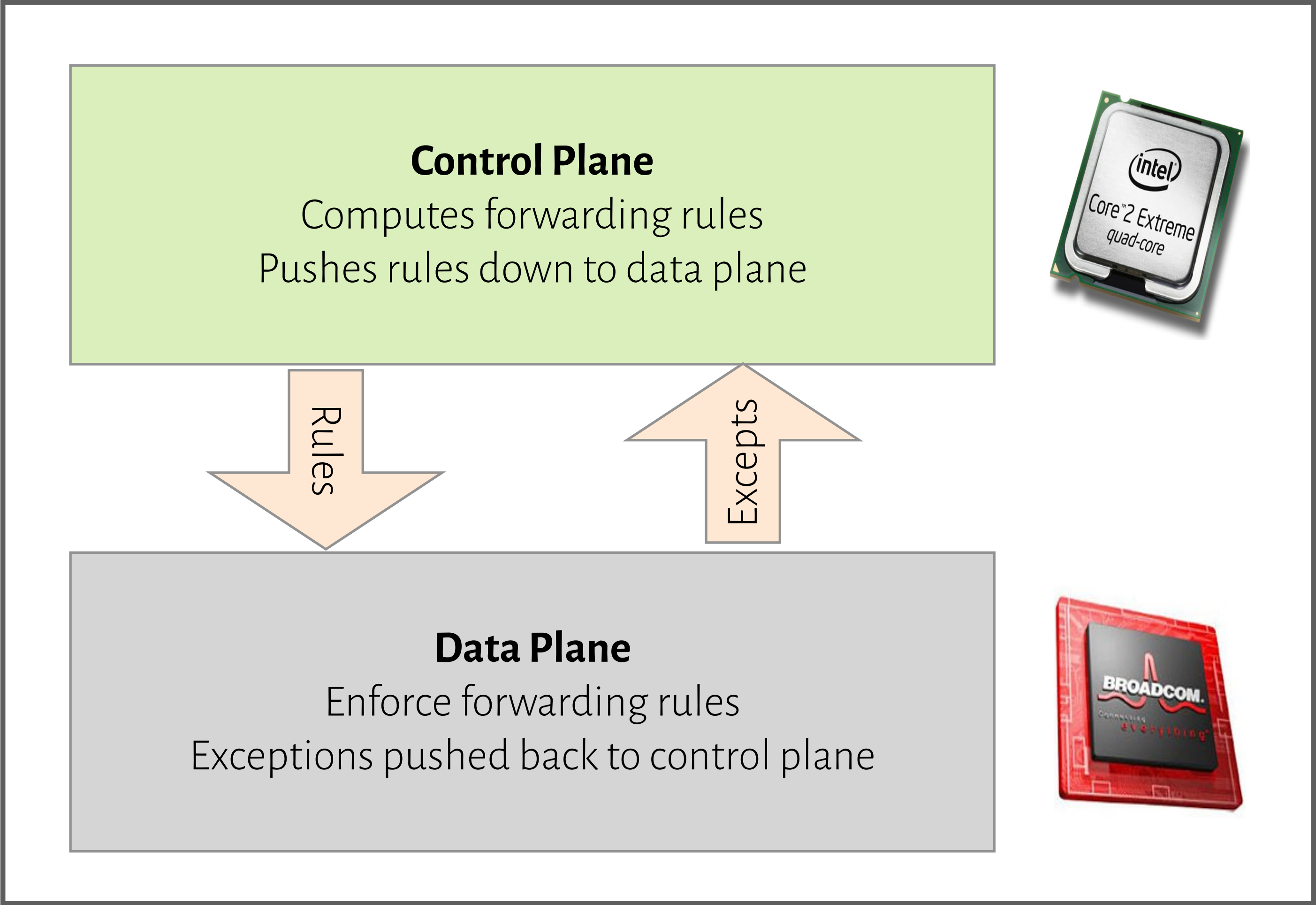


Slicing a traditional network

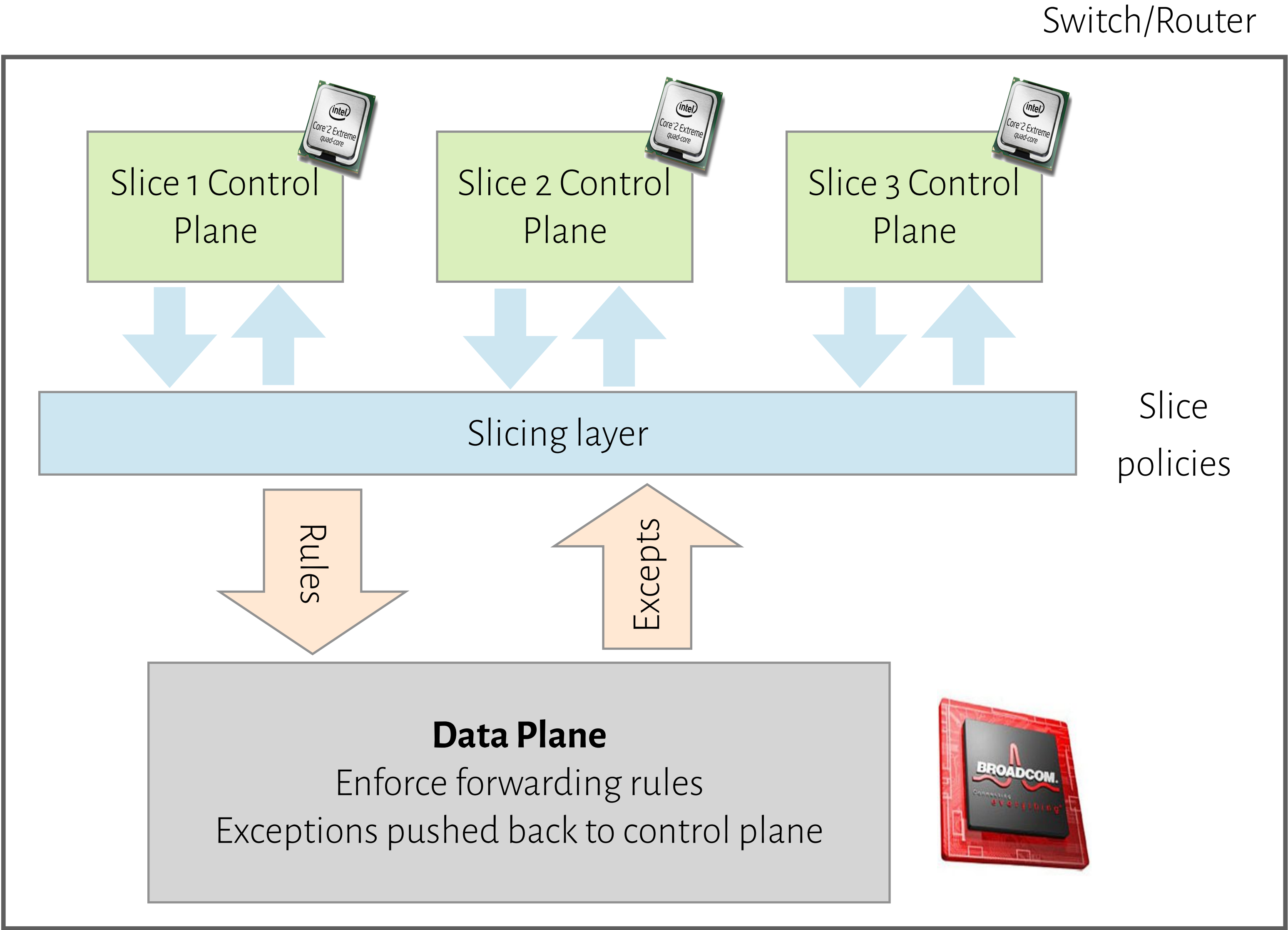


Current network devices

Switch/Router



Slicing layer



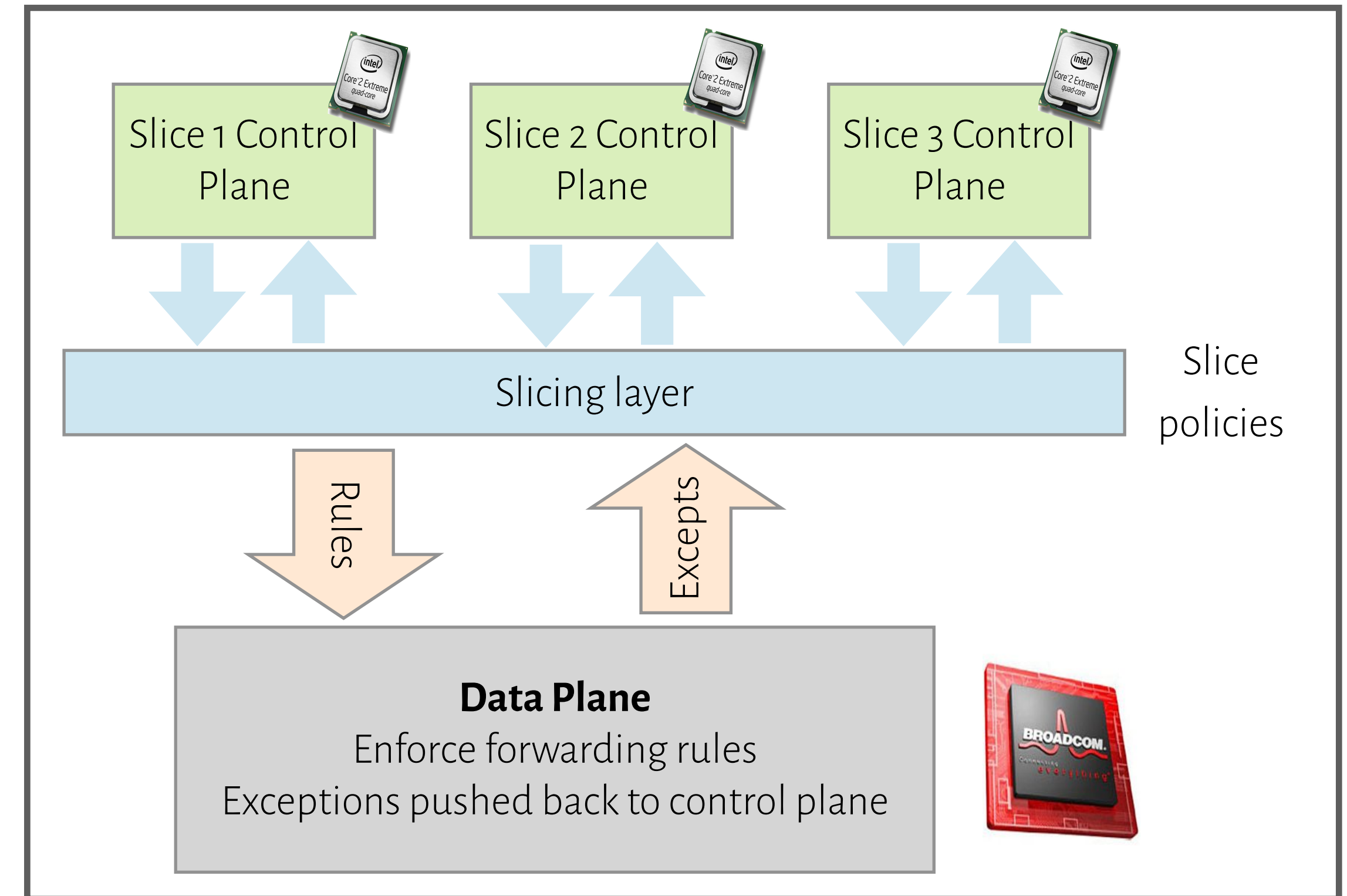
Network slicing architecture

A network slice is a collection of sliced switches/routers

- Data plane is unmodified
- Packets forward with **no performance penalty**
- Slicing with existing ASIC

Transparent slicing layer

- Each slice believes it owns the data path
- Enforces isolation between slides: rewrites, drops rules to adhere to slice policy
- Forwards exceptions to correct slice(s)



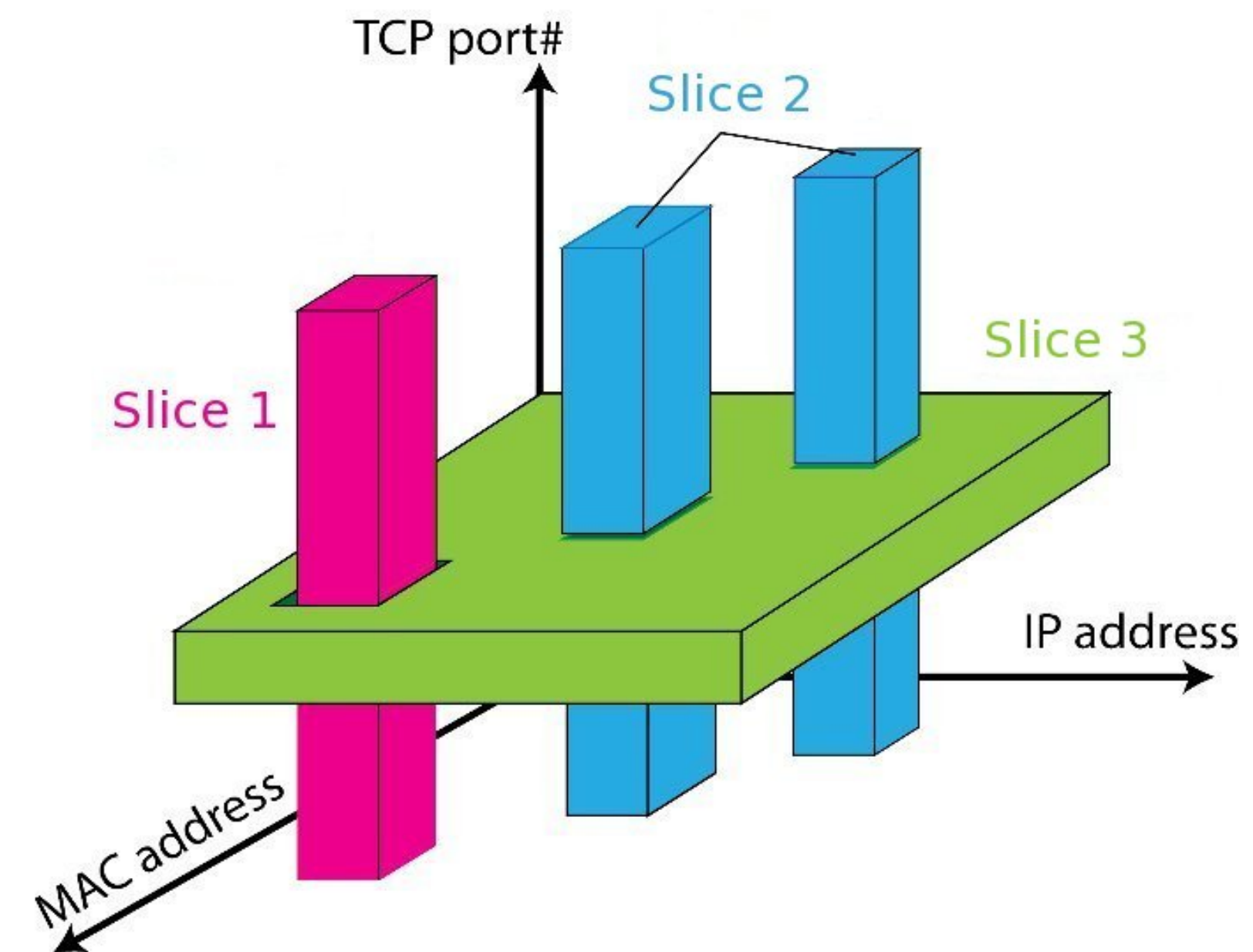
Slicing policies

The slicing policy specifies the **resource limit** for each slice:

- Link bandwidth
- Maximum number of forwarding rules (on switches)
- Topology
- Fraction of switch/router CPU

FlowSpace: which packet does the slice control?

- Maps packets to slices according to their "classes" defined by the packet header fields



Real user traffic: opt-in

Allow users to opt-in to services in real time

- Users can delegate control of individual flows to slices
- Add new FlowSpace to each slice's policy

Examples

- "Slice 1 will handle my HTTP traffic"
- "Slice 2 will handle my VoIP traffic"
- "Slice 3 will handle everything else"

Creates incentives for building high-quality services!



Source: gacovinolack.com

Slice definition

Bob's experimental slice: all HTTP traffic to/from users who opted in

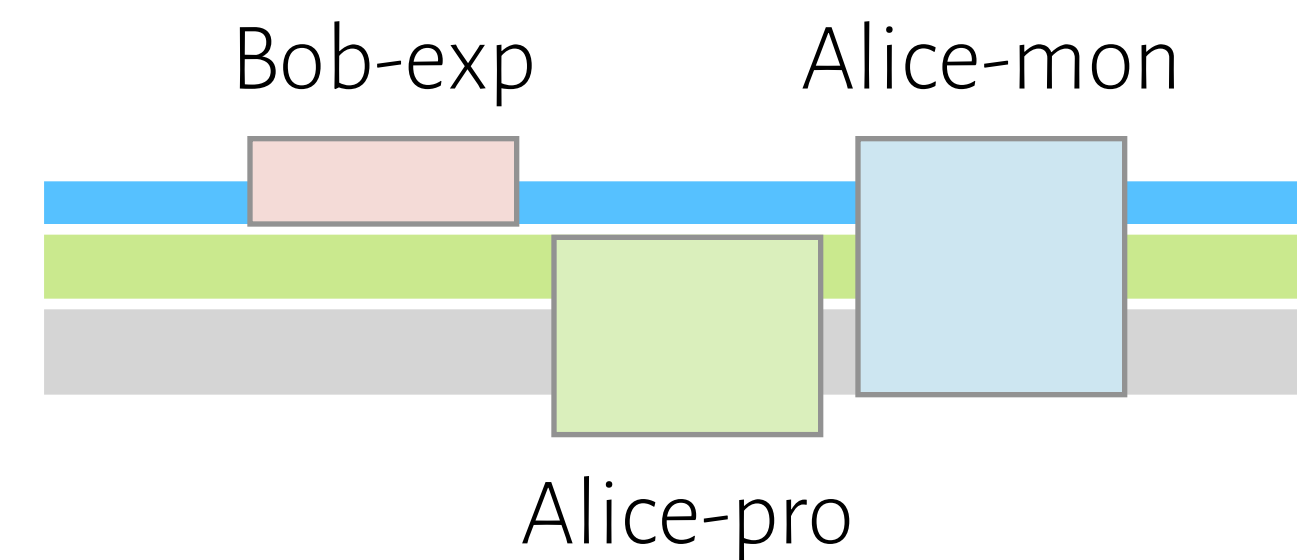
- Allow: tcp_port=80 and ip=user_ip

Alice's production slice: complementary to Bob's slice

- Deny: tcp_port=80 and ip=user_ip
- Allow: all

Alice's monitoring slice: all traffic in all slices

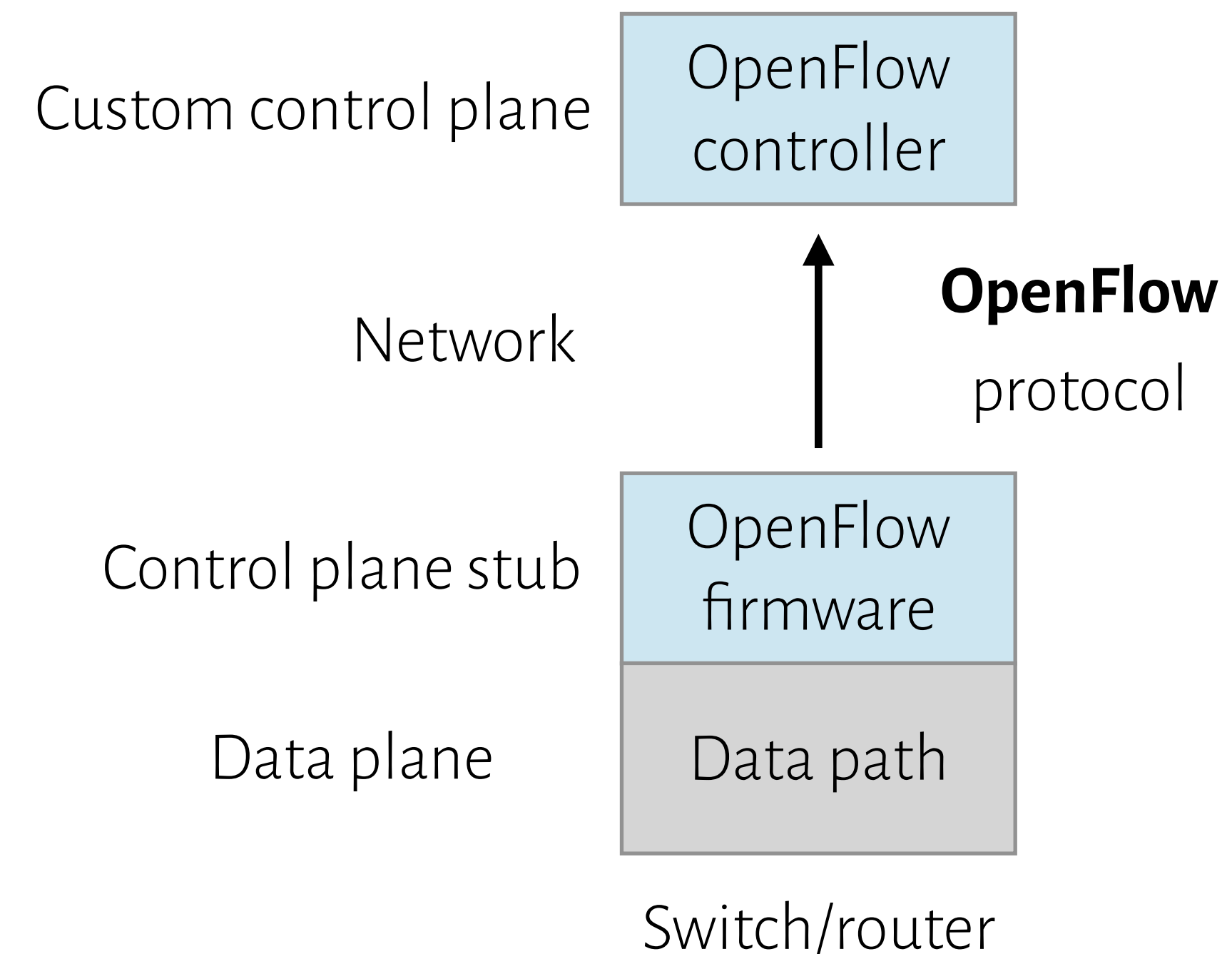
- Read-only: all



Slicing with OpenFlow

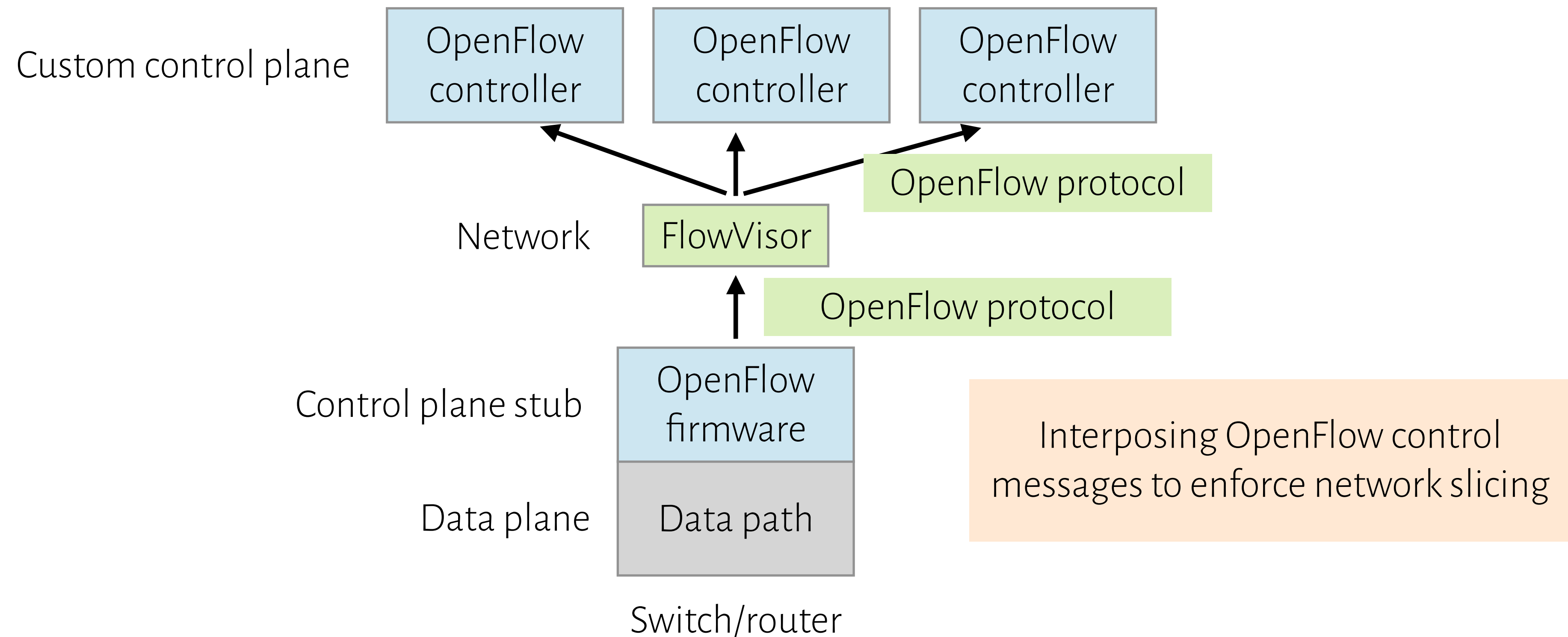
Recall OpenFlow:

- API for controlling packet forwarding
- Abstraction of control/data plane protocols
- Works on commodity hardware (via firmware upgrade)

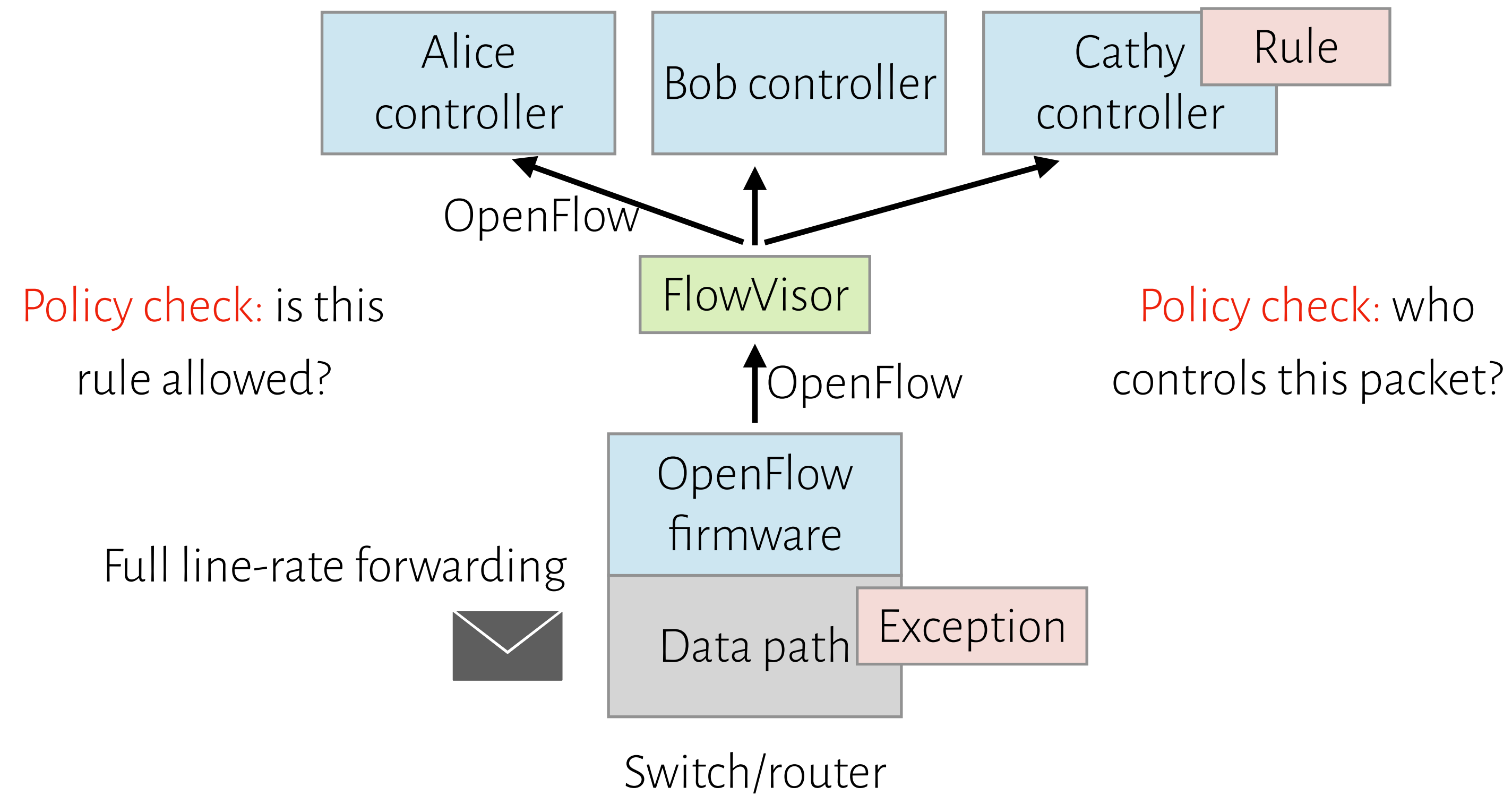


How should we slice an OpenFlow-based software defined network?

FlowVisor

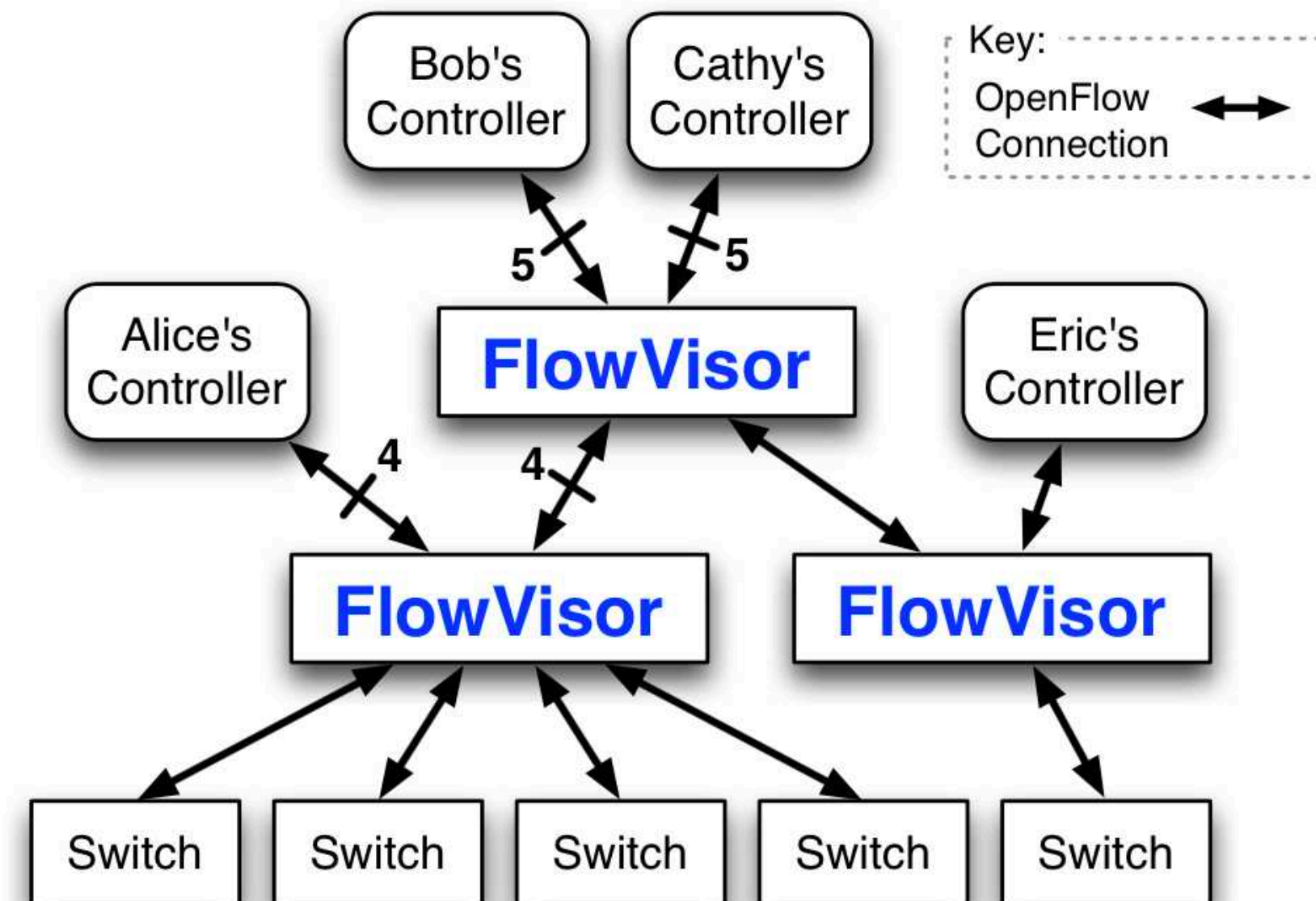


FlowVisor packet handling



Recursive slicing

FlowVisor can trivially recursively slice an already sliced network, creating hierarchies of FlowVisors

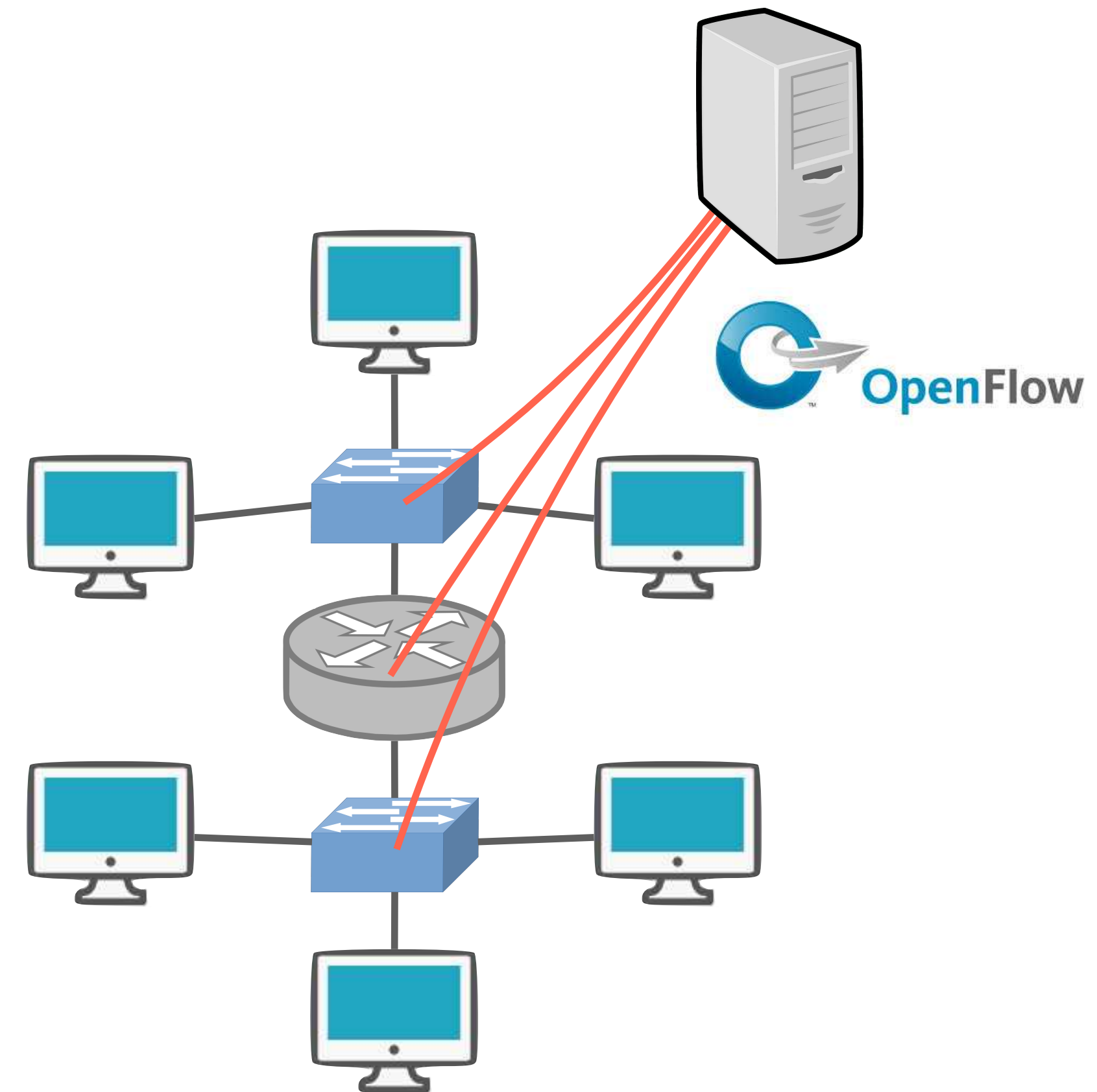


Questions?

Summary

Lecture 8: Software defined networking

- Complexities in network control
- Software defined networking idea
- Abstractions in SDN
- Network virtualization, FlowVisor
- Mininet for experimentation



Next week: programmable data plane

Limitations of OpenFlow

- Bounded to existing protocols
- New protocols/packet formats not supported

How to support arbitrary packet format and enable a fully programmable network?

